





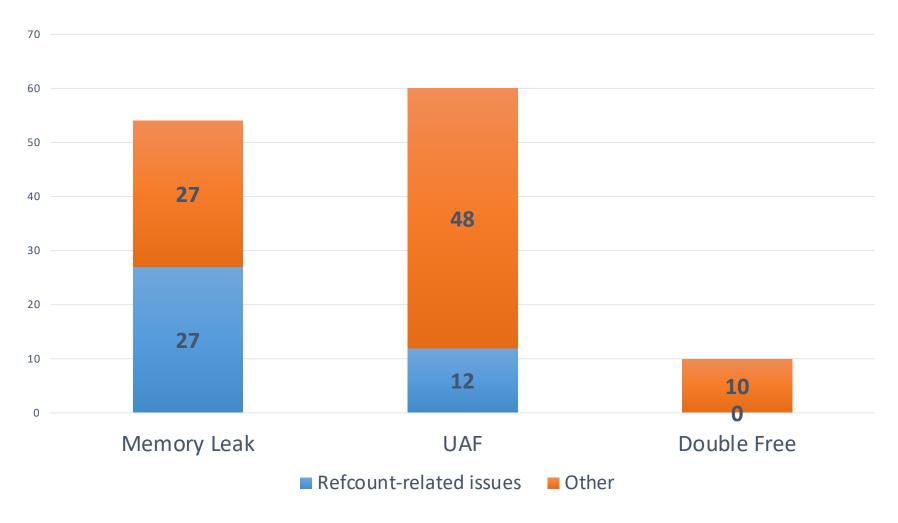
# LinKRID: Vetting Imbalance Reference Counting in Linux Kernel with Symbolic Execution

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# Kernel Memory Management Bugs



A statistic on memory management-related issues in Linux 4.14

#### Reference Counter in Kernels

Key idea: each object includes a ref\_count

```
- Assume obj * p = NULL;
- p = obj1; // obj1->ref_count++
- p = obj2; // obj1->ref_count--, obj2->ref_count++
```

- If an object's ref\_count == 0, it is garbage
  - No pointers target that object
  - thus, it can be safely freed
- Manually managed → error prone

#### Refcount Bugs

- Two types of bugs
  - Failure to increase the counter when a new reference is created -> dangling pointer(s) -> use-after-free (UAF)
  - Failure to decrement the counter when a reference has been removed -> memory leaks

```
-2030,7 +2030,6 @@ static int replace_map_fd_with_map_ptr
                    if (IS_ERR(map)) {
                            verbose ("fd %d is not pointing
                                   insn->imm):
                            fdput(f);
                            return PTR_ERR(map);
               CVE-2016-4557
@@ -794,6 +794,7 @@ long join_session_keyring(const
                 ret = PTR_ERR(keyring);
                 goto error2:
        } else if (keyring == new->session keyring)
                 key_put (keyring);
                 ret = 0:
                 goto error2;
               CVE-2014-0728
```

### Challenge-1: Refcount Bug Modeling

- Existing stratigies to model refcount bugs
  - 1. RID<sup>[1]</sup> and CID<sup>[2]</sup>: IPP-based (Inconsistent Path Pair) modeling
    - Check the inconsistency between path pair which have same arguments and return value
    - Limited checking scenario, e.g., inconsistency within a single path (CVE-2016-0728)
  - 2. Pungi<sup>[3]</sup>: Affine abstraction modeling
    - Check refcount changes equal to the number of escaped references
    - Not special to Linux Kernel, e.g., Linux refcount wrappers
    - High False Positives Rate in Kernel

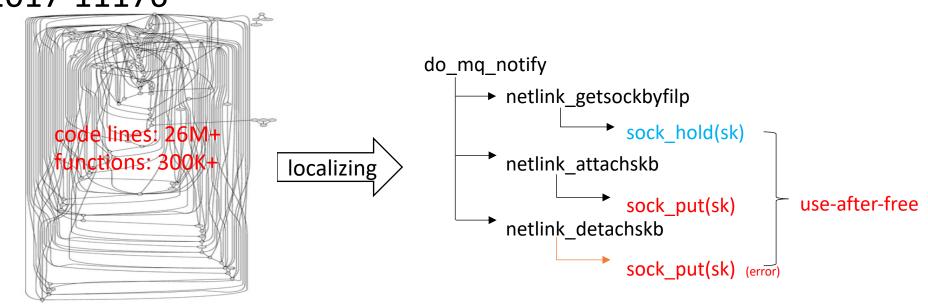
<sup>[1]</sup> Junjie Mao, Yu Chen, Qixue Xiao, and Yuanchun Shi. RID: Finding Reference Count Bugs with Inconsistent Path Pair Checking. ASPLOS 2016.

<sup>[2]</sup> Xin Tan, et al. "Detecting Kernel Refcount Bugs with {Two-Dimensional} Consistency Checking. USENIX Security 2021.

<sup>[3]</sup> Siliang Li and Gang Tan. Finding reference-counting errors in Python/C programs with affine analysis. ECOOP 2014.

### Challenge-2: Refcount Bug Checking

How to scale up the path-sensitive analyse Linux kernel codes, eg.
 CVE-2017-11176



- 1. RID<sup>[1]</sup>: symbolic execution
  - without details on how to scale up
- 2. CID<sup>[2]</sup> and Pungi<sup>[3]</sup>: static analysis

#### Challenge-3: Kernel Refcount Conventions

• Existing works without considering kernel refcount conventions<sup>[1]</sup> (e.g., CVE-2016-4805)

- 1. External reference:
  - strictly abide the reference consistency invariant
- 2. Internal reference (or weak reference):
  - It simply means the object "exists"
  - Violates the refcount consistency invariant
  - Used inside a software cache, such as a radix tree, a double linked list...

```
static void amp_mgr_destroy(struct kref *kref) {
    struct amp_mgr *mgr = container_of(kref, struct amp_mgr, kref);
    mutex_lock(&amp_mgr_list_lock);
    list_del(&mgr->list);
    mutex_unlock(&amp_mgr_list_lock);
    kfree(mgr);
}

int amp_mgr_put(struct amp_mgr *mgr) {
    return kref_put(&mgr->kref,
        &amp_mgr_destroy);
}
```

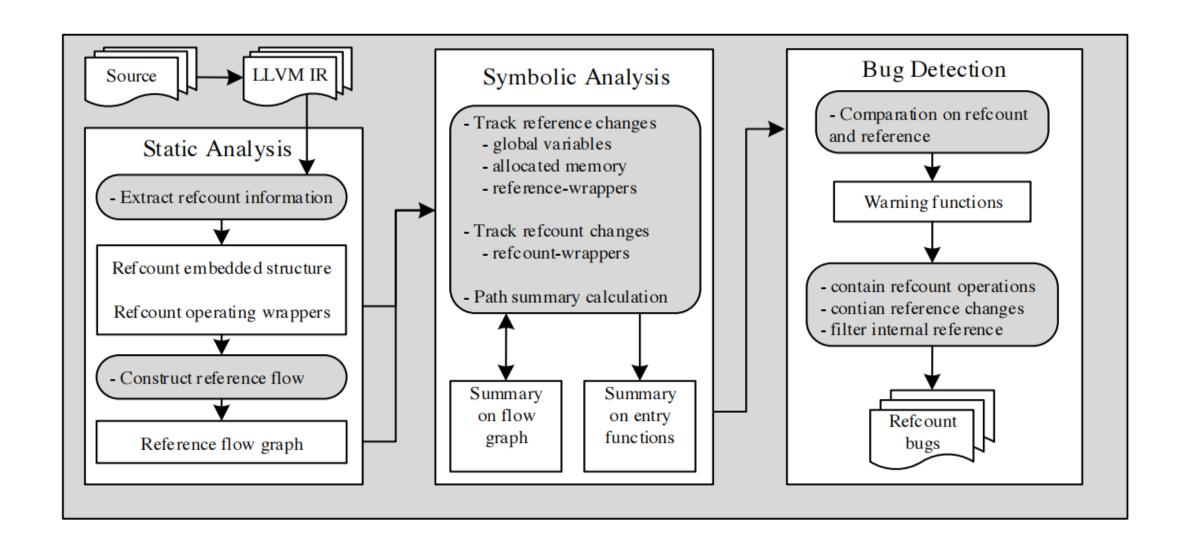
An example of internal reference

# Our Key Observations

- Bug Oracle: the refcount and corresponding reference changes should be consistent
- Invariant: Δ#(reference) == #escaped #released

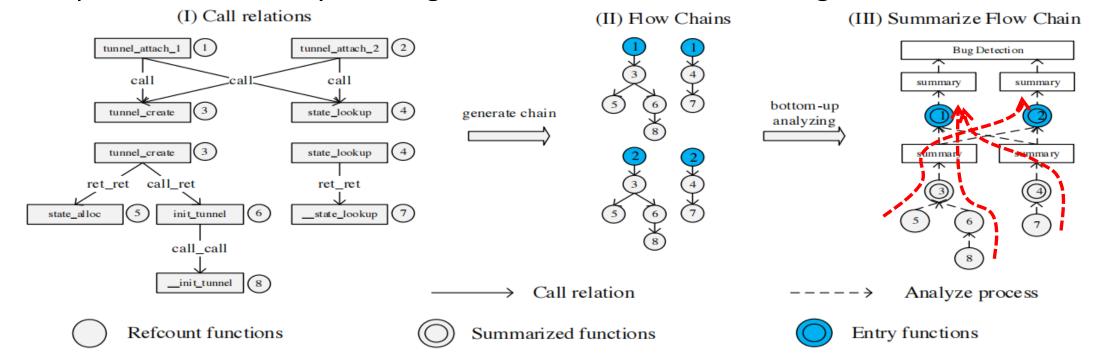
```
func_2(){ // local reference scope begins
   func_1(){ // local reference scope begins
                                                                          p = list_lookup("name", list);
     p = kmalloc(sizeof( *p ));
                                                                          refcnt_get(p); // refcnt += 1
     refcnt_init(p); // init refcnt = 1
     list_insert(p, list); // save to heap, escaped
                                                                          use(p);
        reference += 1
                                                                          refcnt_put(p); // refcnt -= 1
5 } // local reference scope ends
                                                                   14 } // local reference scope ends
  // #escape = 1, #release = 0, \Deltarefcnt = 1.
                                                                        // \#escape = 0, \#release = 0, \Deltarefcnt = 0.
  // \Deltarefcnt == (#escape - #release) == 1. correct!
                                                                        // \Deltarefcnt == (#escape - #release) == 0. correct!
                 (a) Creation
                                                                                         (b) Usage
                                                                       func_4(){ // local reference scope begins
    func_3(){ // local reference scope begins
                                                                          p = list_lookup("name", list);
      p = list_lookup("name", list);
                                                                          list_remove(p, list); // released reference += 1
     refcnt_get(p); // refcnt += 1
                                                                          refcnt_put(p); // refcnt -= 1
      insert(p, list2); // escaped reference += 1
                                                                   30 } // local reference scope ends
22 } // local reference scope ends
                                                                   31 // #escape = 0, #release = 1, \Deltarefcnt = -1.
    // \#escape = 1, \#release = 0, \Deltarefcnt = 1.
    // \Deltarefcnt == (#escape - #release) == 1. correct!
                                                                   32 // \Deltarefcnt == (#escape - #release) == -1. correct!
                                                                                          (d) Release
                 (c) Escape
```

#### LinkRID Workflow



#### Summary-based Chain Analysis

- Ideas: using under-constrained symbolic execution to analyze flow chains symbolically to identify potential buggy paths
  - scalable: bottom-up summary-based computing
  - precious: on-the-fly tracking reference and refcount changes



# Bug Detection (1)

• Given a flow chain, LinkRID reports a potential bug whenever there exists a local reference | r satisfying the following condition

$$\Delta refcount_{/r} \neq |escape_{/r}| - |release_{/r}|$$

- Detailed bug reporting
  - branch direction, refcount get/put, source code line...

```
join_session_keyring path record:
Error Path:
security/keys/process_keys.c L780, get keyring refcnt +
  1 by find_keyring_by_name
security/keys/process_keys.c L781, if condition is false
security/keys/process_keys.c L791, if condition is false
security/keys/process_keys.c L794, if condition is true
...
Summary: <keyring, escape=Ø, release=Ø, Δrefcnt=1, return=0>
Result : refcnt error in keyring
```

# Bug Detection (2)

- Identifying Internal References: two internal reference patterns
  - (1) separate/additional references (released automatically when the refcount reaches zero)
  - (2) released in domain-specific manners, eg.

Considering as Internal References

- Association of Refcount and Reference Changes: heuristic-based aliasing inferring
  - (1) if two local references (e.g., gdm and tty\_port) happen to trigger two warnings in the same function
  - (2) one local reference is "embedded" in another

Considering as related refcount and reference changes

#### Implementation: LinkRID

- Static anlaysis (multiple LLVM Passes with version 9.0.0)
  - Call graph construction from KENALI
  - Data flow analysis
- Under-constraint symoblic execution engine (based on KLEE 2.0)
  - Lazy initialization of Kernel Objects
  - Summary-based symbolic execution
- Bug detector
  - Convention filters
- Totoal: 3.5k lines of C++ code and 2.7k line of python code

#### **Evaluation Setup**

- Target: Linux 4.14.0 compiled with LLVM
  - git commit bebc608
  - Compiled with allyesconfig
  - LLVM 3.9 with make defconfig and make allyesconfig
- Experimental configuration
  - (Virtual machines each with Intel Xeon E3-1220 CPU and 32GB RAM) X 2

#### Performance

- Time to analyze the whole kernel (with allyesconfig)
  - Static analysis: 1 hours
    - Identify 445 refcounted structures, 685 refcounted functions and 54731 related functions
    - generate 12075 flow chains
  - Symoblic execution: 192 hours
    - gnerate 9419 summaries

Phase	Description	D <del>ata</del>	
		defconfig	allyesconfig
Static Analysis	refcounted structures	85	445
	refcount wrapper	149	685
	related functions	16155	54731
	flow chains	4063	12075
Symbolic Execution	summaries	2964	9419
Bug Detection	BUGs	31	118

#### Effectiveness

- Found 118 new refount bugs in Linux Kernel 4.14.0
  - 82 are new bugs, including 30 1-day bugs and 52 0-day bugs.
  - 20 patches have already been merged into the Linux upstream.

#### Conclusion

- LinkRID, a scalable and practical refcount bug discovery tool
  - scale up the path-sensitive analysis, with a lightweight static analysis and an under-constrained symbolic execution
  - a set of conventions seemingly erroneous but benign usage patterns

# Thank you for listening! QA