

Painting Squares in $\Delta^2 - 1$ Shades

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Abstract

Cranston and Kim conjectured that if G is a connected graph with maximum degree Δ and G is not a Moore Graph, then $\chi_\ell(G^2) \leq \Delta^2 - 1$; here χ_ℓ is the list chromatic number. We prove their conjecture; in fact, we show that this upper bound holds even for online list chromatic number.

1 Introduction

Graph coloring has a long history of upper bounds on a graph's chromatic number χ in terms of its maximum degree Δ . A greedy coloring (in any order) gives the trivial upper bound $\chi \leq \Delta + 1$. In 1941, Brooks [4] proved the following strengthening: If G is a graph with maximum degree $\Delta \geq 3$ and clique number $\omega \leq \Delta$, then $\chi \leq \Delta$. In 1977, Borodin and Kostochka [3] conjectured the following further strengthening.

Conjecture 1 (Borodin-Kostochka Conjecture [3]). If G is a graph with $\Delta \geq 9$ and $\omega \leq \Delta - 1$, then $\chi \leq \Delta - 1$.

If true, this conjecture is best possible in two senses. First, the condition $\Delta \geq 9$ cannot be dropped (or even weakened), as shown by the following graph (see Figure 1). Let D_i induce a triangle for each $i \in \{1, \dots, 5\}$; if $|i - j| \equiv 1 \pmod{5}$, then add all edges between vertices of D_i and D_j . This yields an 8-regular graph on 15 vertices with clique number 6 and chromatic number 8; it would be a counterexample to the conjecture if we weakened the hypothesis $\Delta \geq 9$. Similarly, even if we require $\omega \leq \Delta - 2$, we cannot conclude that $\chi \leq \Delta - 2$, as is shown by the join of a clique and a 5-cycle. For each $\Delta \in \{3, \dots, 8\}$, examples are known [6, 13] where $\omega \leq \Delta - 1$ but $\chi = \Delta$. Kostochka has informed us that already in 1977 when he and Borodin posed Conjecture 1, they believed the following stronger “list version” was true; however they omitted this version from their paper, and it appeared in print [6] only in 2013. We define the list chromatic number, denoted χ_ℓ , in Section 2 below.

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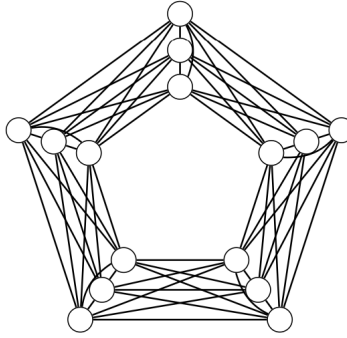


Figure 1: The hypothesis $\Delta \geq 9$ in the Borodin–Kostochka Conjecture is best possible.

Conjecture 2 (Borodin-Kostochka Conjecture (list version)). If G is a graph with $\Delta \geq 9$ and $\omega \leq \Delta - 1$, then $\chi_\ell \leq \Delta - 1$.

The purpose of this paper is to prove the following conjecture of Cranston and Kim [5]. In fact, we will prove this conjecture in the more general setting of online list coloring. It is easy to show, as we do below, that Conjecture 2 implies Conjecture 3.

Conjecture 3 (Cranston-Kim [5]). If G is a connected graph with maximum degree $\Delta \geq 3$, and G is not a Moore graph, then $\chi_\ell(G^2) \leq \Delta^2 - 1$.

A Moore graph is a Δ -regular graph G on $\Delta^2 + 1$ vertices such that $G^2 = K_{\Delta^2+1}$; the sole example when $\Delta = 3$ is the Petersen graph. Hoffman and Singleton [12] famously proved that Moore graphs exist only when $\Delta \in \{2, 3, 7, 57\}$. When $\Delta \in \{2, 3, 7\}$ Moore graphs exist and are known to be unique, and when $\Delta = 57$ no Moore graph is known.

In 2008 Cranston and Kim [5] proved Conjecture 3 when $\Delta = 3$, and suggested that a similar but more detailed approach might prove the whole conjecture. As mentioned above, it is easy to show that Conjecture 3 is implied by Conjecture 2. The key is the following easy lemma at the end of [5]: If G is connected and is not a Moore graph and G has maximum degree $\Delta \geq 3$, then G^2 has clique number at most $\Delta^2 - 1$. The proof is short once we have a result of Erdős, Fajtlowicz, and Hoffman [11] stating that a “near-Moore graph”, i.e., a Δ -regular graph such that $G^2 = K_{\Delta^2}$, exists only when $\Delta = 2$. For details, see the start of the proof of the Main Theorem.

We note that recently Conjecture 3 was generalized to higher powers. Let M denote the maximum possible degree when a graph of maximum degree k is raised to the d th power, i.e., vertices are adjacent in G^d if they are distance at most d in G . Miao and Fan [14] conjectured that if G is connected and G^d is not K_{M+1} , then we can save one color over the bound given by Brooks Theorem, i.e., $\chi(G^d) \leq M - 1$. This was proved by Bonamy and Bousquet [2] in the more general context of online list coloring.

The following conjecture is due to Wegner [20], in the 1970’s. It is a less well-known variant of Wegner’s analogous conjecture when the class \mathcal{G}_k is restricted to planar graphs.

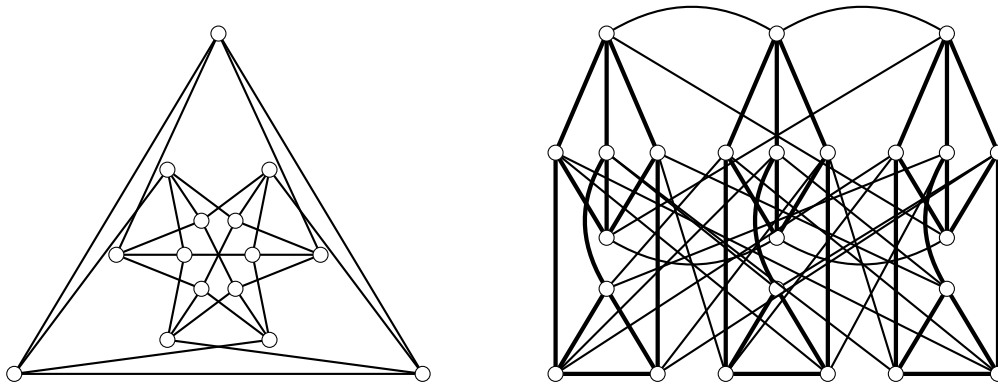


Figure 2: On the left is a 4-regular graph G_1 such that $G_1^2 = K_{15}$.
 On the right is a 5-regular graph G_2 such that $G_2^2 = K_{24}$.

Conjecture 4 (Wegner [20]). For each fixed k , let \mathcal{G}_k denote the class of all graphs with maximum degree at most k and form \mathcal{G}_k^2 by taking the square G^2 of each graph G in \mathcal{G}_k . Now $\max_{H \in \mathcal{G}_k^2} \chi(H) = \max_{H \in \mathcal{G}_k^2} \omega(H)$.

Wegner in fact posed a more general conjecture for all powers of \mathcal{G}_k ; however, here we restrict our attention to Conjecture 4, specifically for small values of k . For each $H \in \mathcal{G}_k^2$, we have $\Delta(H) \leq k^2$, so Brooks' Theorem implies that $\chi(H) \leq k^2$ unless some component of H is K_{k^2+1} . For $k = 1$ Wegner's Conjecture is trivial. For $k \in \{2, 3, 7\}$ it is easy; in each case \mathcal{G}_k contains a Moore graph G , and letting $H = G^2$, we have $H = K_{k^2+1}$, so $\chi(H) = \omega(H) = k^2 + 1$. Thus, the first two open cases of Conjecture 4 are $k = 4$ and $k = 5$. Our Main Theorem shows that every graph G in \mathcal{G}_4 satisfies $\chi_\ell(G^2) \leq 15$ and every graph G in \mathcal{G}_5 satisfies $\chi_\ell(G^2) \leq 24$. Matching lower bounds are shown in Figure 2: we have $G_1 \in \mathcal{G}_4$ with $\omega(G_1^2) = 15$ and $G_2 \in \mathcal{G}_5$ with $\omega(G_2^2) = 24$. Both graphs were discovered by Elspas ([9] and p. 14 of [15]) and are known to be the unique graphs G with $\Delta \in \{4, 5\}$ and $G^2 = K_{\Delta^2-1}$. This confirms Wegner's Conjecture when $k = 4$ and $k = 5$.

Rather than coloring, or even list coloring, this paper is about *online list coloring*, a generalization introduced in 2009 by Schauz [16] and Zhu [22], and the *online list chromatic number*, χ_p , also called the *paint number*. We give the definition in Section 2, but for now if you are unfamiliar with χ_p , you can substitute χ_ℓ (or even χ) and the Main Theorem remains true. Our main result is the following.

Main Theorem. *If G is a connected graph with maximum degree $\Delta \geq 3$ and G is not the Peterson graph, the Hoffman-Singleton graph, or a Moore graph with $\Delta = 57$, then $\chi_p(G^2) \leq \Delta^2 - 1$.*

We conclude this section with the following conjecture, which generalizes our Main Theorem as well as Conjecture 2.

Conjecture 5 (Borodin-Kostochka Conjecture (online list coloring version)). *If G is a graph with $\Delta \geq 9$ and $\omega \leq \Delta - 1$, then $\chi_p \leq \Delta - 1$.*

The structure of the paper is as follows. In Section 2 we give background and definitions. In Section 3, we prove the Main Theorem, subject to a number of lemmas about forbidden subgraphs in a minimal counterexample. In Section 4 we prove the lemmas that we deferred in Section 3. Finally, in Section 5, we generalize the online list chromatic number to the Alon–Tarsi number, and extend our Main Theorem to that setting.

2 Preliminaries

Here we give definitions and background. Most of our terminology and notation is standard. We write $A \setminus B$ for $A \cap \overline{B}$. If H is a subgraph of G , then $G \setminus H$ means $G[V(G) \setminus V(H)]$, that is G with the vertices of H deleted. For graphs G and H , the *join* $G \vee H$ is formed from the disjoint union of G and H by adding all edges with one endpoint in each of $V(G)$ and $V(H)$. For any undefined terms, see West [21].

A *list size assignment* $f : V(G) \rightarrow \mathbb{Z}^+$ assigns to each vertex in G a list size. An *f-assignment* L assigns to each vertex v a subset of the positive integers $L(v)$ with $|L(v)| = f(v)$. An L -coloring is a proper coloring ϕ such that $\phi(v) \in L(v)$ for all v . A graph G is *f-list colorable* (or *f-choosable*) if G has an L -coloring for every f -assignment L . In particular, we are interested in the case where $f(v) = k$ for all v and some constant k . The *list chromatic number* of G or *choice number* of G , denoted $\chi_\ell(G)$, or simply χ_ℓ when G is clear from context, is the minimum k such that G is k -choosable. List coloring was introduced by Vizing [19] and Erdős, Rubin, and Taylor [10] in the 1970s. Both groups proved the following extension of Brooks' Theorem. If G is a graph with maximum degree $\Delta \geq 3$ and clique number $\omega \leq \Delta$, then $\chi_\ell \leq \Delta$.

The next idea we need came about 30 years later. In 2009, Schauz [16] and Zhu [22] independently introduced the notion of online list coloring. This is a variation of list coloring in which the list sizes are determined (each vertex v gets $f(v)$ colors), but the lists themselves are provided online by an adversary.

We consider a game between two players, *Lister* and *Painter*. In round 1, Lister presents the set of all vertices whose lists contain color 1. Painter must then use color 1 on some independent subset of these vertices, and cannot change this set in the future. In each subsequent round k , Lister chooses some subset of the uncolored vertices to contain color k in their lists, and Painter chooses some independent subset of these vertices to receive color k . Painter wins if he succeeds in painting all vertices. Alternatively, Lister wins if he includes a vertex v among those presented on each of $f(v)$ rounds, but Painter never paints v .

A graph is *online k-list colorable* (or *k-paintable*) if Painter can win whenever $f(v) = k$ for all v . The minimum k such that a graph G is online k -list colorable is its *online list chromatic number*, or *paint number*, denoted χ_p . A graph is *d₁-paintable* if it is paintable when $f(v) = d(v) - 1$ for each vertex v . In [7], the authors introduced *d₁-choosable* graphs, which are the list-coloring analogue. Interest in d_1 -paintable graphs owes to the fact that none can be induced subgraphs of a minimal graph with maximum degree Δ that is not $(\Delta - 1)$ -paintable. In particular, if G is a minimal counterexample to our Main Theorem, then G^2 contains no induced d_1 -paintable subgraph.

Lemma 1. *Let G be a graph with maximum degree Δ and H be an induced subgraph of G that is d_1 -paintable. If $G \setminus H$ is $(\Delta - 1)$ -paintable, then G is $(\Delta - 1)$ -paintable.*

Proof. Let G and H satisfy the hypotheses. We give an algorithm for Painter to win the online coloring game when $f(v) = \Delta - 1$ for all v . Painter will simulate playing two games simultaneously: a game on $G \setminus H$ with $f(v) = \Delta - 1$ and a game on H with $f(v) = d_H(v) - 1$. Let S_k denote the set of vertices presented by Lister on round k . Painter first plays round k of the game on $G \setminus H$, pretending that Lister listed the vertices $S_k \setminus H$. Let I_k denote the independent set of these that Painter chooses to color k .

Let $S'_k = (S_k \cap V(H)) \setminus N(I_k)$, the vertices of H that are in S_k and have no neighbor in I_k . Now Painter plays round k of the game on H , pretending that Lister listed S'_k . Each vertex in $V(G \setminus H)$ will clearly be listed $\Delta - 1$ times. Consider a vertex v in $V(H)$. It will appear in $N(I_k)$ for at most $d_G(v) - d_H(v)$ rounds. So v will appear in S'_k for at least $(\Delta - 1) - (d_G(v) - d_H(v)) \geq d_H(v) - 1$ rounds. Now Painter will win both simulated games, and thus win the actual game on G . \square

When the graph G in Lemma 1 is a square, we immediately get that $G \setminus H$ is $(\Delta - 1)$ -paintable, as we note in the next lemma.

Lemma 2. *Let G be a connected graph with maximum degree Δ and let H be an induced subgraph of G^2 . If H is d_1 -paintable, then G^2 is d_1 -paintable. If there exists v with $d_{G^2}(v) < \Delta^2 - 1$, then G^2 is $(\Delta^2 - 1)$ -paintable.*

Proof. We prove the first statement first. Let $V = V(G)$ and $V_1 = V(H)$. Clearly a graph is d_1 -paintable only if each component is. So we assume that $G^2[V_1]$ is connected. For simplicity, we assume also that $G[V_1]$ is connected. If not, then some vertex v has neighbors in two or more components of $G[V_1]$. We simply add v to V_1 , since we can color v first (when it still has at least two uncolored neighbors).

Form G' from G by contracting $G[V_1]$ to a single vertex r . Let T be a spanning tree in G' rooted at r . Let σ be an ordering of the vertices of $G \setminus H$ by nonincreasing distance in T from r . Each time that Lister presents a list of vertices, Painter chooses a maximal independent subset of them, by greedily adding vertices in order σ . Each vertex $v \in V \setminus V_1$ is followed in σ by the first two vertices on a path in T from v to r . Thus v will be colored. We now combine strategies for $G^2 \setminus H$ and H as in the proof of Lemma 1.

Now we prove the second statement, which has a similar proof. Suppose there exists v with $d_{G^2}(v) < \Delta^2 - 1$. As before we order the vertices by nonincreasing distance in some spanning tree T from v , and we put v and some neighbor u last in σ . The difference now is that even for u and v we are given $\Delta^2 - 1$ colors. Since $d_{G^2}(v) < \Delta^2 - 1$, either (i) v lies on a 3-cycle or 4-cycle or else (ii) $d_G(v) < \Delta$ or v has some neighbor u with $d_G(u) < \Delta$; in Case (ii), by symmetry we assume $d_G(v) < \Delta$. In Case (i), $d_{G^2}(u) \leq \Delta^2 - 1$ for some neighbor u of v on the short cycle and by assumption $d_{G^2}(v) < \Delta^2 - 1$; so the two final vertices of σ are u and v . In Case (ii), we again have $d_{G^2}(v) < \Delta^2 - 1$ and $d_{G^2}(u) \leq \Delta^2 - 1$, so again u and v are last in σ . \square

The previous lemma implies that $\Delta^2 - 1 \leq d_{G^2}(v) \leq \Delta^2$ for every vertex v in a graph G such that G^2 is not $(\Delta^2 - 1)$ -paintable. A vertex v is *high* if $d_{G^2}(v) = \Delta^2$, and otherwise

it is *low*. The proof of Lemma 2 proves something slightly more general, which we record in the following corollary.

Corollary 3. *Let G be a graph with maximum degree Δ and let H be an induced subgraph of G^2 . Let $f(v) = d(v) - 1$ for each high vertex of G^2 and $f(v) = d(v)$ for each low vertex. If H is f -paintable, then G^2 is $(\Delta^2 - 1)$ -paintable.*

Now we will introduce the Alon–Tarsi Theorem, but we need a few definitions first. Let G be a graph and let \vec{D} be a digraph arising by orienting the edges of G . A *circulation* is a subgraph of \vec{D} in which each vertex has equal indegree and outdegree; circulations are also called eulerian subgraphs. The parity of a circulation is the parity of its number of edges. For a digraph \vec{D} , let $EE(\vec{D})$ (resp. $EO(\vec{D})$) denote the set of circulations that are even (resp. odd).

Theorem A (Alon and Tarsi [1]). *For a digraph \vec{D} , if $|EE(\vec{D})| \neq |EO(\vec{D})|$, then \vec{D} is f -choosable, where $f(v) = 1 + d_{\vec{D}}(v)$ for all v .*

The proof that Alon and Tarsi gave was algebraic and not constructive. In their paper, they asked for a combinatorial proof. This was provided by Schauz [17], in the more general setting of paintability. His proof relies on an elaborate inductive argument. The argument does yield a constructive algorithm, although in general it may run in exponential time. In [18], Schauz proved an online version of the combinatorial nullstellensatz from which the paintability version of Alon and Tarsi’s theorem can also be derived.

Theorem B (Schauz [17]). *For a digraph \vec{D} , if $|EE(\vec{D})| \neq |EO(\vec{D})|$, then \vec{D} is f -paintable, where $f(v) = 1 + d_{\vec{D}}(v)$ for all v .*

Our main result relies heavily on forbidding d_1 -paintable subgraphs. For many of the smaller d_1 -paintable graphs that we need, we give direct proofs. However, for some of the larger d_1 -paintable graphs, particularly the classes of unbounded size, our proofs of d_1 -paintability use Theorem B.

3 Proof of Main Theorem

In this section we prove our main result, subject to a number of lemmas on forbidden subgraphs, which we defer to the next section. We typically prove that a subgraph is forbidden by showing that it is d_1 -paintable. If a copy of a subgraph H in G^2 contains low vertices, then this configuration is reducible as long as H is f -paintable, where $f(v) = d_H(v) - 1$ for each high vertex v and $f(w) = d_H(w)$ for each low vertex w . For many of the graphs, we give an explicit winning strategy for Painter. In contrast, for some of the graphs, particularly those of unbounded size, we don’t give explicit winning strategies. Instead, we show that they are d_1 -paintable via Schauz’s extension of the Alon–Tarsi Theorem (Theorem B).

Main Theorem. *If G is a connected graph with maximum degree $\Delta \geq 3$ and G is not the Peterson graph, the Hoffman–Singleton graph, or a Moore graph with $\Delta = 57$, then $\chi_p(G^2) \leq \Delta^2 - 1$.*

Proof. Let G be a connected graph with maximum degree $\Delta \geq 3$, other than the graphs excluded in the Main Theorem. Assume that G^2 is not $(\Delta^2 - 1)$ -paintable. By Lemma 2, if there exists $v \in V(G)$ with $d_{G^2}(v) < \Delta^2 - 1$, then G^2 is $(\Delta^2 - 1)$ -paintable. So G is Δ -regular and has girth at least 4. Further, no vertex of G lies on two or more 4-cycles. It will be helpful in what follows to show that $\omega(G^2) \leq \Delta^2 - 1$.

Clearly $\Delta(G^2) \leq \Delta^2$. Further, $\omega(G^2) = \Delta^2 + 1$ only if $G^2 = K_{\Delta^2+1}$. Hoffman and Singleton [12] showed this is possible only if $\Delta \in \{2, 3, 7, 57\}$; such a graph G is called a Moore graph. When $\Delta \in \{2, 3, 7\}$, the unique realizations are the 5-cycle, the Peterson graph, and the Hoffman-Singleton graph. When $\Delta = 57$, no realization is known. These are precisely the graphs excluded from the theorem. Now we consider the case $\omega(G^2) = \Delta^2$. Erdős, Fajtlowicz, and Hoffman [11] showed that the only graph H such that $H^2 = K_{\Delta(H)^2}$ is C_4 . Cranston and Kim noted that if H^2 is not a clique on at least Δ^2 vertices, then in fact $\omega(H^2) \leq \Delta^2 - 1$. For completeness, we reproduce the details.

Suppose that $\omega(G^2) = \Delta^2$, and let U be the vertices of a maximum clique in G^2 . The result of Erdős, Fajtlowicz, and Hoffman implies that U is not all of V . Choose $v, w \in V$ with $v \in U$, $w \notin U$ and v adjacent to w . Since $d_{G^2}(v) = \Delta^2$ and $w \notin U$, every neighbor of w must be in U . Applying the same logic to these neighbors, every vertex within distance 2 of w must be in U . But now we can add w to U to get a larger clique in G^2 . This contradiction implies that in fact $\omega(G^2) \leq \Delta^2 - 1$.

Two vertices are *linked* if they are adjacent in G^2 , and otherwise they are *unlinked*. When we write that vertices are adjacent or nonadjacent, we mean in G ; otherwise we write linked or unlinked. We write $v \leftrightarrow w$ if v and w are adjacent, and $v \not\leftrightarrow w$ otherwise.

Now we are ready to present the details of the proof. Before that, it is useful to give a general outline. Our approach is to show that G^2 must contain a forbidden induced subgraph. Above, we noted that G must be Δ -regular and have girth at least 4. Now we consider the possibilities for the girth of G , which we denote by g . Suppose $g = 4$. For each vertex v on a 4-cycle, $d_{G^2}(v) \leq \Delta^2 - 1$. So it is straightforward to show that G contains a forbidden subgraph.

When G has sufficiently high girth, the situation is also simpler than the general case. Now we let U denote the vertices of a shortest cycle, as well as a few off-cycle neighbors. Because the girth is high enough, we know that $G^2[U]$ has only the edges guaranteed by its definition. In other words, no pairs of vertices in U have edges in G^2 due to common neighbors outside of U ; if they did, then we could find a shorter cycle. This allows us again to show that G^2 contains a forbidden subgraph. Each of our forbidden subgraphs contains an induced cycle in G , so the case of high girth is interesting, since we need to show that an infinite family of subgraphs (with unbounded girth) are forbidden. It turns out the the girth is sufficiently high for this approach to succeed when $g \geq 9$. This leaves us with the cases $g \in \{5, 6, 7, 8\}$.

The case $g = 6$ is quite easy, since C_6^2 is a forbidden subgraph. The cases $g = 7$ and $g = 8$ are not hard, but involves some tedious details. So the real difficulty is in the case $g = 5$. This makes sense, since the Moore graphs have girth 5. At some point, we must

explicitly use our assumption that G is not a Moore graph. This argument requires some detailed structural analysis of G . It is this structural analysis that makes up much of the proof in this section. Most of the other possibilities for girth are deferred to lemmas that we prove in the next section.

Case 1: G has girth 4

Let C be a 4-cycle with vertices v_1, \dots, v_4 , and let $\mathcal{C} = V(C)$. It is helpful to note that every v_i is low. We need two lemmas. These were first proved in [8] for list coloring, and we generalize them to online list coloring in Lemmas 5 and 6. The following two configurations in G^2 are reducible: (A) $K_4 \vee \overline{K_2}$ where some vertex $w \in V(K_4)$ is low and (B) $K_3 \vee \overline{K_2}$ where some vertices $w \in V(K_3)$ and $x \in V(\overline{K_2})$ are both low.

Note that $G^2[\mathcal{C}] \cong K_4$. This implies that every w adjacent to some $v_i \in \mathcal{C}$ must be linked to all of \mathcal{C} . Suppose not, and let w be adjacent to v_1 and not linked to v_3 . Now $G^2[\mathcal{C} \cup \{w\}] \cong K_3 \vee \overline{K_2}$, and every v_i is low; this is (B), which is forbidden. Now suppose that w_1 and w_2 are vertices adjacent to v_i and v_j , respectively. We must have w_1 linked to w_2 , since otherwise $G^2[\mathcal{C} \cup \{w_1, w_2\}]$ is (A), which is forbidden.

Now let x be a vertex at distance 2 from v_1 and not adjacent to any v_i ; let w_1 be a common neighbor of v_1 and x . Since w_1 is linked to v_3 , they have a common neighbor w_3 . Now x is linked to v_1, w_1 , and w_3 . To avoid configuration (B), x must be linked to all of \mathcal{C} . Thus, all vertices within distance 2 of v_1 must be linked to all of \mathcal{C} . Now every pair of vertices x and y that are both within distance 2 of v_1 must be linked; otherwise $G^2[\mathcal{C} \cup \{x, y\}]$ is (A). So the vertices within distance 2 of v_1 induce in G^2 a clique of size Δ^2 , which contradicts that $\omega(G^2) \leq \Delta^2 - 1$.

Case 2: G has girth at least 5

Let g denote the girth of G . First suppose that $g = 6$, and let U be the vertices of a 6-cycle. Note that $G^2[U] = C_6^2$, since girth 6 implies there are no extra edges. Since C_6^2 is d_1 -paintable, by Lemma 9, we are done by Lemma 2.

Suppose $g = 7$. Let U denote the vertices of some 7-cycle in G , with a pendant edge at a single vertex of the cycle. Because G has girth 7, $G^2[U]$ has only the edges guaranteed by its definition. We show in Lemma 18 that $G^2[U]$ is d_1 -paintable. So again, we are done by Lemma 2.

Suppose instead that $g \geq 8$. Let $U = \{v_1, \dots, v_g, w_1, w_5\}$ be the vertices of some g -cycle in G together with pendant edges v_1w_1 and v_5w_5 . If $g \geq 9$, then $G^2[U]$ has only the edges guaranteed by its definition. If $g = 8$, then $G^2[U]$ has the edges guaranteed by its definition as well as possibly the extra edge w_1w_5 . For each girth g at least 8, we show in Lemma 19 and Lemma 16 that $G^2[U]$ is d_1 -paintable. So again, we are done by Lemma 2.

Now we consider girth 5. Our approach is similar to that for girth 4, but we must work harder since we don't necessarily have any low vertices. Let C be a 5-cycle with vertices v_1, \dots, v_5 . Let $k = \Delta - 2$. For each i , let V_i denote the neighbors of v_i not on C . Let $\mathcal{C} = V(C)$ and let $\mathcal{D} = \cup_{i=1}^5 V_i$. Each vertex of \mathcal{D} is linked to either 5, 4, or 3 vertices of \mathcal{C} . We call these B_0 -vertices, B_1 -vertices, and B_2 -vertices, respectively (a B_i -vertex is unlinked to i vertices of \mathcal{C}). We will consider four possibilities for the number and location of each type of vertex. In each case we find a d_1 -paintable subgraph. Let L denote the

subgraph $G[\mathcal{D}]$. Since G has girth 5, we have $\Delta(L) \leq 2$. Each vertex w with $d_L(w) = 2 - i$ is a B_i -vertex (for $i \in \{0, 1, 2\}$).

Suppose that G has two B_1 -vertices w_1 and w_2 and they are unlinked with distinct vertices in \mathcal{C} . Let $H = G^2[\mathcal{C} \cup \{w_1, w_2\}]$. If w_1 and w_2 are linked, then $H = K_3 \vee C_4 \supset K_2 \vee C_4$, which is d_1 -paintable, by Lemma 10. If instead w_1 and w_2 are unlinked, then $H = K_3 \vee P_4$, which is also d_1 -paintable, by Lemma 11. So we assume that all B_1 -vertices are unlinked with the same vertex $v \in \mathcal{C}$. As a result, each B_1 -vertex is an endpoint of a path of length 3 (mod 5) in L , for otherwise the two endpoints of the path are unlinked with different vertices in \mathcal{C} (here we use that G contains no 3-cycle and no 4-cycle). Since the number of odd degree vertices in any graph is even, here the number of B_1 -vertices is even.

Case 2.1: G has a B_1 -vertex w_1 and a B_2 -vertex w_2 .

Let $H = G^2[\mathcal{C} \cup \{w_1, w_2\}]$. Suppose the four vertices of \mathcal{C} linked to w_1 include the three vertices of \mathcal{C} linked to w_2 . If w_1 and w_2 are linked, then $H = K_3 \vee P_4$, and if w_1 and w_2 are unlinked, then $H = K_3 \vee (K_1 + P_3)$. In each case, H is d_1 -paintable, by Lemmas 11 and 12, respectively.

Suppose instead that the four vertices of \mathcal{C} linked to w_1 do not include all three vertices of \mathcal{C} linked to w_2 . If w_1 is linked with w_2 , then $H \supset K_2 \vee C_4$, which is d_1 -paintable by Lemma 10. If w_1 is unlinked with w_2 , then we will finish by Lemma 15, but we need a little explanation first.

Since each B_1 -vertex is an endpoint of a path in L with length 3 (mod 5), if G has any B_1 -vertices, then it has B_1 -vertices adjacent to two successive cycle vertices. So, if G has a B_1 -vertex and a B_2 -vertex, then we may choose them so that their corresponding cycle vertices are either identical or nonadjacent. If these cycle vertices are identical, then the B_2 -vertex is linked with three cycle vertices contained among those linked with the B_1 -vertex. So the situation is as in the first paragraph of Case 2.1. If these cycle vertices are nonadjacent, then the situation is covered by Lemma 15, so H is again d_1 -paintable. Thus, G^2 cannot contain both B_1 -vertices and B_2 -vertices.

Case 2.2: G has no B_1 -vertices, but only some B_2 -vertices, and possibly also B_0 -vertices.

Now L consists of disjoint cycles, each with length a multiple of 5. This implies that each V_i contains the same number of B_2 -vertices; by assumption this number is at least 1. We call a pair of B_2 vertices with distinct cycle neighbors *near* if their cycle neighbors are adjacent and *far* if their cycle neighbors are nonadjacent. If any pair of far B_2 -vertices are linked, then G has a d_1 -paintable subgraph, by Lemma 13. If any pair of near B_2 -vertices are linked, then, together with their adjacent cycle vertices, they induce $K_2 \vee C_4$, which is d_1 -paintable by Lemma 10. Thus, we consider the subgraph induced by \mathcal{C} and 3 non-successive B_2 -vertices, say with cycle neighbors v_1, v_2, v_4 . Each such subgraph is d_1 -paintable, by Lemma 14.

Combining Cases 2.1 and 2.2, we conclude that G contains no B_2 -vertices.

Case 2.3: G has B_1 -vertices (and also B_0 -vertices).

Recall that G has an even number of B_1 -vertices and they are all unlinked with the same vertex. By symmetry, assume that G has B_1 -vertices $w_2 \in V_2$ and $w_3 \in V_3$ and they

are both unlinked with v_5 . We will find two disjoint pairs of nonadjacent vertices, such that all four are linked with $\mathcal{C} - v_5$.

Since w_3 is a B_1 -vertex, it is the endpoint of some path in L ; let $w_1 \in V_1$ be the neighbor of w_3 on this path. We will show that w_1 is unlinked with some vertex in \mathcal{D} .

Recall that $|\mathcal{D}| = 5k$. Suppose that w_1 is linked to each vertex of \mathcal{D} . Since $d_L(w_1) = 2$ and $d_L(w_3) = 1$, at most 3 of these $5k - 1$ vertices linked with w_1 can be reached from w_1 by following edges in L . Clearly w_1 is linked to the other $k - 1$ vertices of V_1 . Now for each vertex w of the remaining $(5k - 1) - 3 - (k - 1) = 4k - 3$ vertices in \mathcal{D} , w_1 must have a common neighbor x with w and $x \notin \mathcal{D} \cup \mathcal{C}$. Furthermore, each such common neighbor x can link w_1 to at most 4 of these vertices (at most one in each other V_i , since the girth is 5). However, this requires at least $\lceil \frac{4k-3}{4} \rceil = k$ additional neighbors of w_1 , but we have already accounted for 3 neighbors of w_1 . Thus, w_1 is unlinked with some vertex $y \in \mathcal{D}$.

Let z be a B_1 -vertex distinct from y , e.g., either w_2 or w_3 will do. Now z and v_5 are unlinked and w_1 and y are unlinked. But every vertex of $\{w_1, v_5, y, z\}$ is linked to $\mathcal{C} - v_5$. Thus $G^2[(\mathcal{C} - v_5) \cup \{w_1, v_5, y, z\}] = K_4 \vee H$, where H contains disjoint pairs of nonadjacent vertices. So $K_4 \vee H$ is d_1 -paintable, by Lemma 7.

Case 2.4: \mathcal{D} has only B_0 -vertices.

Let $H = G^2[\mathcal{C} \cup \mathcal{D}]$. We will show that if H is not a clique, then we can choose a different 5-cycle and be in an earlier case. Suppose that H is not a clique. Since \mathcal{D} is linked to \mathcal{C} and $G^2[\mathcal{C}] = K_5$, we must have $w_1, w_2 \in \mathcal{D}$ with w_1 and w_2 unlinked. By symmetry, we have only two cases.

First suppose that $w_1 \in V_1$ and $w_2 \in V_2$ and w_1 and w_2 are unlinked. Since w_1 is a B_0 -vertex, we have $w_3 \in V_3$ with $w_1 \leftrightarrow w_3$. Consider the 5-cycle $w_1 v_1 v_2 v_3 w_3$. Now w_2 is not linked to w_1 , which makes w_2 not a B_0 -vertex for that 5-cycle. So we are in Case 2.1, 2.2, or 2.3 above. Now suppose instead that $w_1 \in V_1$ and $w_3 \in V_3$ and w_1 and w_3 are unlinked. Now we pick some $w'_3 \in V_3$ with $w_1 \leftrightarrow w'_3$ and consider the 5-cycle $w_1 v_1 v_2 v_3 w'_3$. Since w_3 and w_1 are unlinked, w_3 is not a B_0 -vertex for this 5-cycle, so we are in Case 2.1, 2.2, or 2.3 above. Hence $G^2[\mathcal{C} \cup \mathcal{D}]$ must be a clique.

To link all vertices in \mathcal{D} , we must have $k(k - 1)$ additional vertices in G , at distance 2 from \mathcal{C} ; call the set of them \mathcal{F} . We see that $|\mathcal{F}| \geq k(k - 1)$ as follows. All $\binom{5k}{2}$ pairs of vertices in \mathcal{D} are linked. The $5\binom{k}{2}$ pairs contained within a common V_i are linked via vertices of \mathcal{C} . Each of the $5k$ vertices is linked with exactly 4 vertices via edges of L . The remaining links all must be due to vertices of \mathcal{F} , and each vertex of \mathcal{F} can link at most $\binom{5}{2} = 10$ pairs of vertices in \mathcal{D} (at most one vertex in each V_i , since G has girth 5). Thus $|\mathcal{F}| \geq ((\binom{5k}{2}) - 5\binom{k}{2} - 5k(4)/2) / \binom{5}{2} = k(k - 1)$. If any vertex $x \in \mathcal{F}$ has fewer than exactly one neighbor in each V_i , then some pair of vertices in \mathcal{D} will be unlinked. Thus, each $x \in \mathcal{F}$ has exactly one neighbor in each V_i . This implies that \mathcal{F} is linked to \mathcal{C} , and hence that $|\mathcal{F}| = k(k - 1)$. We will show that every pair of vertices in $\mathcal{C} \cup \mathcal{D} \cup \mathcal{F}$ is linked.

Suppose there exists $w \in \mathcal{D}$ and $x \in \mathcal{F}$ with w and x unlinked. By symmetry, we assume $w \in V_1$. There exist $w_1 \in V_1$ and $w_2 \in V_2$ with $x \leftrightarrow w_1$ and $x \leftrightarrow w_2$. Now consider the 5-cycle $xw_1 v_1 v_2 w_2$. Since w and x are unlinked, w is not a B_0 -vertex for that 5-cycle. This puts us in Case 2.1, 2.2, or 2.3 above. So \mathcal{F} must be linked to \mathcal{D} .

Finally suppose there exist $x_1, x_2 \in \mathcal{F}$ with x_1 and x_2 unlinked. Now there exist

$w_1, w_2 \in V_1$ with $x_1 \leftrightarrow w_1$ and $x_2 \leftrightarrow w_2$. Since G has girth 5, we have $x_1 \not\leftrightarrow w_2$. And since x_1 is linked with w_2 , they have some common neighbor $y \in \mathcal{D} \cup \mathcal{F}$. Now consider the 5-cycle $x_1 w_1 v_1 w_2 y$. Since x_1 and x_2 are unlinked, x_2 is not a B_0 -vertex for this 5-cycle. Hence, we are in Case 2.1, 2.2, or 2.3.

Thus, all vertices of $\mathcal{C} \cup \mathcal{D} \cup \mathcal{F}$ are pairwise linked. Now $|\mathcal{C} \cup \mathcal{D} \cup \mathcal{F}| = 5 + 5k + k(k-1) = k^2 + 4k + 5 = (k+2)^2 + 1 = \Delta^2 + 1$. This contradicts that $\omega(G^2) \leq \Delta^2 - 1$ and completes the proof. \square

We note that many of the cases of the above proof actually prove that G^2 is d_1 -paintable (rather than only $(\Delta(G)^2 - 1)$ -paintable), and hence has paint number at most $\Delta(G^2) - 1$. In particular, this is true when G has girth 6, 7, or at least 9. Probably with more work, we could also adapt the proof to the case when G has girth 8. The Conjecture that G^2 is $(\Delta(G^2) - 1)$ -paintable unless $\omega(G^2) \geq \Delta(G^2)$ is a special case of Conjecture 5. The main obstacle to proving this stronger result is the case when G has girth at most 5, particularly girth 3 or girth 4.

4 Proofs of forbidden subgraph lemmas

In what follows, we slightly abuse the terminology of high and low vertices defined earlier. Now a vertex is *high* if its list size is one less than its degree and *low* if its list size equals its degree. Note that if a vertex v is high (resp. low) in G by our old definition, then it will be high (resp. low) in each induced subgraph H by our new definition. A vertex is *very low* if its list size is greater than its degree. When a vertex v in a graph G is very low, we may say that we *delete* v . If $G - v$ is paintable from its lists, then so is G . On each round, we play the game on $G - v$ and consider v after all other vertices, coloring it only if its list contained the color for that round and we have colored none of its neighbors on that round. Recall that S_k denotes the vertices with lists containing color k . We write E_k for the empty graph on k vertices, i.e., $E_k = \overline{K_k}$. In what follows, all vertices not specified to be low are assumed to be high.

4.1 Direct proofs

For pictures of the graphs in Lemmas 4 through 12, see Figures 6 and 6 in Section 5.

Lemma 4. *If G is $K_4 - e$ with one degree 3 vertex high and the other vertices low, then G is f -paintable.*

Proof. Let v_1, v_2 denote the degree 3 vertices, with v_1 low, and let w_1, w_2 denote the degree 2 vertices. If $w_1, w_2 \in S_1$, then color them both with 1. Now the remaining vertices are low and very low, so we can finish. Otherwise, color some v_i with 1, choosing v_2 if possible. Now at least one w_j becomes very low and the uncolored v_k is low, so we can finish. \square

Lemma 5. *If G is $K_3 \vee E_2$ with a low vertex in the K_3 and a low vertex in the E_2 , then G is f -paintable.*

Proof. Denote the vertices of the K_3 by v_1, v_2, v_3 , with v_1 low, and the vertices of E_2 by w_1, w_2 , with w_1 low. If $w_1, w_2 \in S_1$, then color them both 1. Now v_1 becomes very low and v_2 and v_3 each become low, so we finish greedily, ending with v_2 and v_1 . Suppose $w_2 \in S_1$. If $v_2 \in S_1$ (or $v_3 \in S_1$, by symmetry), then color v_2 with 1. Now w_1 becomes very low (since $S_1 \not\supseteq \{w_1, w_2\}$), and v_1 remains low, so we can finish greedily. If instead $v_1 \in S_1$ and $v_2, v_3 \notin S_1$, then color v_1 with 1. Again w_1 becomes very low and v_2 and v_3 become low, so we can finish greedily. The situation is similar if S_1 contains only a single w_i . Thus, $w_2 \notin S_1$. Since $S_1 \neq \{w_1\}$, some v_i is in S_1 . Use color 1 on v_i , choosing v_2 or v_3 if possible. What remains is $K_4 - e$ with one degree 3 vertex high and all others low (or very low). So we finish by Lemma 4. \square

Lemma 6. *If G is $K_4 \vee E_2$ with a low vertex in the K_4 , then G is f -paintable.*

Proof. Denote the vertices of the K_4 by v_1, \dots, v_4 , with v_1 low and the vertices of E_2 by w_1, w_2 . If $w_1, w_2 \in S_1$, then color them both 1. Now v_1 becomes very low and the other v_i become low, so we can finish by coloring greedily, with v_1 last. So S_1 contains at most one w_i , say w_2 . Suppose S_1 contains a v_j other than v_1 . Color v_j with 1. Now w_1 becomes low, v_1 remains low, and the other vertices remain high. So we can finish the coloring by Lemma 5. If the only v_i in S_1 is v_1 , then color it 1. Now the other v_j become low, so again we finish by Lemma 5. Finally, if the only vertex in S_1 is w_2 , then color it 1. Now v_1 becomes very low, and the other v_i become low, so again we can finish by coloring greedily, ending with a low vertex and a very low vertex. \square

Lemma 7. *If G is $K_4 \vee H$ with H containing two disjoint nonadjacent pairs, then G is d_1 -paintable.*

Proof. We may assume $|H| = 4$. Denote the vertices of K_4 by v_1, \dots, v_4 and the vertices of H by w_1, \dots, w_4 with $w_1 \not\leftrightarrow w_2$ and $w_3 \not\leftrightarrow w_4$. If $w_1, w_2 \in S_1$, then color w_1 and w_2 with 1. Now every v_i becomes low, so we can finish by Lemma 6. Similarly, if $w_3, w_4 \in S_1$.

If some v_i is missing from S_1 , then use 1 to color either some v_j or some w_k . In the first case, we finish by Lemma 5 and in the second by Lemma 6. So color v_4 with 1. Now, by symmetry, $w_2, w_4 \notin S_1$, so they each become low. If $w_1, w_2 \in S_2$, then color them both with 2. Now every v_i becomes low, so we can finish by Lemma 5. Similarly if $w_3, w_4 \in S_2$. So S_2 contains at most one of w_1, w_2 and at most one of w_3, w_4 . If S_2 contains no v_i , then we color some w_j with 2. This makes every v_i low. Now we can finish by Lemma 5. So S_2 contains some v_i , say v_3 .

Color v_3 with 1. Recall that S_1 was missing at least one of w_1, w_2 and at least one of w_3, w_4 . (i) If $w_2, w_4 \notin S_2$, then they both become very low, so we can delete them. This in turn makes v_1 and v_2 both very low, so we can finish greedily. (ii) If $w_2, w_3 \notin S_2$, then w_2 becomes very low, so we delete it. Now v_1 and v_2 become low; also w_3 and w_4 are low. Since v_1, v_2, w_3, w_4 induce $K_4 - e$ with all vertices low, we can finish by Lemma 4. By symmetry, this handles the case $w_1, w_4 \notin S_2$. (iii) If $w_1, w_3 \notin S_2$, then the uncolored vertices induce $K_2 \vee H$, with all vertices of H low. Now consider S_3 . If S_3 contains a nonadjacent pair in H , then color them both 3. This makes v_1 and v_2 low, so what remains is $K_4 - e$ with all vertices low. We now finish by Lemma 4. Similarly, if S_3 contains no

v_i , then color some w_j with 3, and we can finish by Lemma 4. So S_3 contains some v_i , say v_2 , and we color v_2 with 3. Now one of w_1, w_2 becomes very low and one of w_3, w_4 becomes very low. We can delete the very low vertices, which in turn makes v_1 very low. We can now finish greedily, since what remains is a 3-vertex path with two low vertices and a very low vertex. \square

We won't use Lemma 8 in the proof, but it is generally useful so we record it here.

Lemma 8. *If G is $K_6 \vee E_3$, then G is d_1 -paintable.*

Proof. Denote the vertices of K_6 by v_1, \dots, v_6 and the vertices of E_3 by w_1, w_2, w_3 . If $w_1, w_2, w_3 \in S_1$, then color w_1, w_2, w_3 all with 1. Now all v_i are very low, so we finish greedily. If no v_i appears in S_1 , then color some w_j with 1. Now all the v_i are low, so we can finish by Lemma 6. So some v_i is in S_1 , say v_6 . Color v_6 with 1. This makes some w_i low, say w_3 . Repeating this argument, we get by symmetry that $v_5 \in S_2$ and S_2 is missing some w_j . If S_2 is missing w_3 , then color v_5 with 2. Now w_3 becomes very low, so we delete it. This in turn makes all uncolored v_k low. Now we can finish by Lemma 6. So instead S_2 is missing (by symmetry) w_2 . Again repeating the argument, we must have $v_4 \in S_3$ and $w_1 \notin S_3$; otherwise we finish by Lemma 5 or Lemma 6. Now we color v_4 with 3. What remains is $K_3 \vee E_3$ with every w_i low.

Now consider S_4 . If $w_1, w_2, w_3 \in S_4$, then color them all with 3. Now all remaining vertices become very low, so we finish greedily. Suppose instead that $w_1 \in S_4$ and $v_1 \notin S_4$. Color w_1 with 4. What remains is $K_3 \vee E_2$ with both w_i low and some v_j low. So we can finish by Lemma 4. A similar approach works for any $w_i \in S_4$ and $v_j \notin S_4$. So instead, assume by symmetry that $v_1 \in S_4$ and $w_1 \notin S_4$. Color v_1 with 4. Now w_1 becomes very low, so we delete it. This in turn makes v_2 and v_3 low. Now we can finish by Lemma 4. \square

Lemma 9. *If G is C_6^2 , then G is d_1 -paintable.*

Proof. Denote the vertices of the 6-cycle by v_1, \dots, v_6 in order. So v_i is adjacent to all but $v_{(i+3) \bmod 6}$. Consider S_1 . If S_1 contains some nonadjacent pair, then color them with 1. What remains is C_4 with all vertices low, so we can complete the coloring since C_4 is 2-paintable. So assume that S_1 contains no nonadjacent pairs. Now without loss of generality, we assume $S_1 = \{v_1, v_2, v_3\}$, since adding vertices to S_1 only makes things harder to color, as long as S_1 induces a clique; we may also need to permute a nonadjacent pair. Color v_1 with 1.

Now v_5 and v_6 become low. Consider S_2 . Again, if S_2 contains a nonadjacent pair, then we color both vertices with 2 and can finish greedily since all remaining vertices are low, except for one that is very low. If $v_2, v_3 \in S_2$, then color v_2 with 2. Now v_6 becomes very low and v_5 remains low, so we can finish greedily. So S_2 misses at least one of v_2, v_3 . Suppose $v_4 \in S_2$. Color v_4 with 3. What remains is C_4 . If $v_2, v_3 \notin S_2$, then all vertices are low, and we can finish since C_4 is 2-paintable. Otherwise, v_5 or v_6 becomes very low and the other remains low. Now we can finish greedily. So $v_4 \notin S_2$. If $v_2 \in S_2$, then color v_2 with 2. Now v_3 and v_4 become low, so we can finish by Lemma 4. An analogous argument works if $v_3 \in S_2$. So assume $v_2, v_3, v_4 \notin S_2$. Now color v_5 or v_6 with 2. Again we can finish by Lemma 4. \square

Lemma 10. *If G is $K_2 \vee C_4$, then G is d_1 -paintable.*

Proof. Denote the vertices of K_2 by v_1, v_2 and the vertices of C_4 by w_1, \dots, w_4 in order. If S_1 contains a pair of nonadjacent vertices, then color them both 1. What remains is $K_4 - e$, with all vertices low. So we can finish by Lemma 4. So S_1 misses at least one of w_1, w_3 and at least one of w_2, w_4 . By symmetry, say it misses w_1 and w_2 . Suppose $v_1, v_2 \notin S_1$. Now by symmetry $w_3 \in S_1$, so color w_3 with 1. This makes each of w_2, v_1, v_2 low. So what remains is $K_3 \vee E_2$ with two low vertices in the K_3 and a low vertex in the E_2 . Hence, we can finish by Lemma 5.

So instead (by symmetry) $v_2 \in S_1$. Color v_2 with 1. What remains is $K_1 \vee C_4$ with w_1 and w_2 low. Consider S_2 . Again if S_2 contains a nonadjacent pair, then we color them both 2, and we can finish greedily. Suppose that $w_3 \in S_2$. If $w_4 \notin S_2$, then we color w_3 with 4; now w_4 becomes low, so we can finish by Lemma 4. If instead $w_4 \in S_2$, then $w_2 \notin S_2$. Now when we color w_3 with 2, w_2 becomes very low, so we can finish greedily. So assume $w_3, w_4 \notin S_2$. If $v_1 \in S_2$, then color v_1 with 1. What remains is C_4 with all vertices low. Now we can finish the coloring since C_4 is 2-paintable. The proof is similar to that for 2-choosability, so we omit it. So assume that $v_1 \notin S_2$. By symmetry, we have $w_1 \in S_2$. Color w_1 with 2. What remains is $K_4 - e$ with only w_3 high. Hence we can finish by Lemma 4. \square

Lemma 11. *If G is $K_3 \vee P_4$, then G is d_1 -paintable.*

Proof. Let v_1, v_2, v_3 denote the vertices of K_3 and w_1, \dots, w_4 denote the vertices of the P_4 in order. If $w_1, w_3 \in S_1$, then color them both 1. Now what remains is $K_3 \vee E_2$ with all but one vertex low, so we can finish by Lemma 5. An analogous strategy works if $w_2, w_4 \in S_1$. So assume S_1 misses at least one of w_1, w_3 and at least one of w_2, w_4 . If S_1 misses v_1 , then use color 1 on some w_j , choosing w_2 or w_3 if possible. Again, we can finish by Lemma 5. So assume $v_1 \in S_1$. Now color v_3 with 1. What remains is $K_2 \vee P_4$ with at least two vertices of the P_4 low. Consider S_2 . If $w_1, w_3 \in S_2$ (or $(w_2, w_4 \in S_2)$), then color them both 2, and we can finish greedily since all vertices are low except for one that is very low. If $v_2 \in S_2$, then color it with 2. Now in each case we can finish by repeatedly deleting very low vertices, possibly using Lemma 4. So $v_2 \notin S_2$ (and by symmetry $v_3 \notin S_2$). If possible use color 2 on w_1 or w_4 . This leaves $K_3 \vee E_2$ with enough low vertices to finish by Lemma 5. Finally, if $w_1, w_4 \notin S_2$, then by symmetry $w_2 \in S_2$, so color w_2 with 2. What remains contains a $K_4 - e$ with all vertices low, so we can finish by Lemma 4. \square

Lemma 12. *If G is $K_3 \vee (K_1 + P_3)$, then G is d_1 -paintable.*

Proof. Let v_1, v_2, v_3 denote the vertices of K_3 ; let w_1, w_2, w_3 denote the vertices of P_3 in order, and let w_4 be the K_1 . If $w_1, w_3 \in S_1$, then color them both 1 and we can finish by Lemma 5. If instead $w_2, w_4 \in S_1$, then color them both 1, and again we can finish by Lemma 5. If $S_1 = \{w_4\}$, then color w_4 with 1. What remains is $K_3 \vee P_3$ with all vertices of the K_3 low. Since $K_3 \vee P_3 \cong K_4 \vee E_2$, we can finish by Lemma 6. If $w_1 \in S_1$ (or $w_2 \in S_1$ or $w_3 \in S_1$) and $v_3 \notin S_1$, then color w_1 with 1. Again we can finish by Lemma 5. This implies that $v_3 \in S_1$.

Since $v_3 \in S_1$, color v_3 with 1. Now at least one of w_1, w_3 becomes low and at least one of w_2, w_4 becomes low. What remains is $K_2 \vee (K_1 + P_3)$, and by symmetry either (i) w_1 and w_2 are low or (ii) w_1 and w_4 are low. Consider (i). If we ignore w_4 , then what remains is $K_2 \vee P_3 \cong K_3 \vee E_2$. Since w_1 and w_2 are low, we can finish by Lemma 5. Instead consider (ii). If $w_1, w_3 \in S_2$, then color them both with 2. What remains is $K_4 - e$ and all vertices are low, so we finish by Lemma 4. Suppose instead that $w_2, w_4 \in S_2$. Color them both with 2, which makes v_1 and v_2 low. If w_1 became very low, then we finish greedily. Otherwise w_3 became low, so we finish by Lemma 4. Now suppose $v_1 \in S_2$, and color v_1 with 2. We have four possibilities. If w_2 and w_3 become low, then we can finish by Lemma 4. Similarly, if w_4 becomes very low, we delete it; now v_2 becomes low, so we can finish by Lemma 4. In the two remaining cases, we can finish greedily by repeatedly deleting very low vertices. \square

4.2 Proofs via the Alon–Tarsi Theorem

Our goal in each of the next lemmas is to prove that a certain graph is d_1 -paintable. Recall the definitions of EE and EO from Section 2, preceding Theorem A. For a digraph \vec{D} , we write $\text{diff}(\vec{D})$ to denote $|EE(\vec{D})| - |EO(\vec{D})|$. In each case we find an orientation \vec{D} such that each vertex has indegree at least 2 and $\text{diff}(\vec{D}) \neq 0$. Now the Alon–Tarsi Theorem, specifically the generalization in Theorem B, proves the graph is d_1 -paintable. To compute $\text{diff}(\vec{D})$, sometimes we simply calculate $|EE(\vec{D})|$ and $|EO(\vec{D})|$ and take the difference. However, this is typically quite difficult if we are considering a set of subgraphs of unbounded size. In that case we want to avoid calculating $|EE(\vec{D})|$ and $|EO(\vec{D})|$ explicitly. Rather, we look for a parity-reversing bijection that pairs elements of $EE(\vec{D})$ with elements of $EO(\vec{D})$. In computing $\text{diff}(\vec{D})$, we can ignore all circulations paired by such a bijection. We also use the following trick to reduce our work. We explain it via an example, but it holds more generally.

Let \vec{D} contain a 5-clique and two other vertices w_1 and w_2 such that for each v either $d^+(v) \leq 3$ or $d^+(v) = 4$ and $w_1, w_2 \in N^+(v)$. In computing $\text{diff}(\vec{D})$, we want to restrict the difference to the set of circulations in which $d^+(w_1) \geq 1$ and $d^+(w_2) \geq 1$; call this $\text{diff}'(\vec{D})$. By inclusion-exclusion, we have $\text{diff}'(\vec{D}) = \text{diff}(\vec{D}) - \text{diff}(\vec{D} - w_1) - \text{diff}(\vec{D} - w_2) + \text{diff}(\vec{D} - w_1 - w_2)$. So it suffices to show that the final three terms on the right side are 0. If any term were nonzero, then, by the Alon–Tarsi Theorem, we would be able to color the corresponding subgraph from lists of size at most 4. However, the subgraph contains a 5-clique, making this impossible. Thus, each term is 0, and we have the desired equality. (In some cases we use a slight variation of this approach, instead concluding that the induced subgraph H with $\text{diff}(H) \neq 0$ is d_1 -paintable.) Finally, we combine this technique with the parity-reversing bijection mentioned above, by restricting the bijection only to the set of circulations where $d^+(w_1) \geq 1$ and $d^+(w_2) \geq 1$.

For brevity, in this section our proofs of Lemmas 13, 14, 15, and 16 consist simply of orientations in which each vertex has in-degree at least 2 and $|EE| \neq |EO|$. In each case, we provide the sizes of EE and EO , which are easy to check with a computer

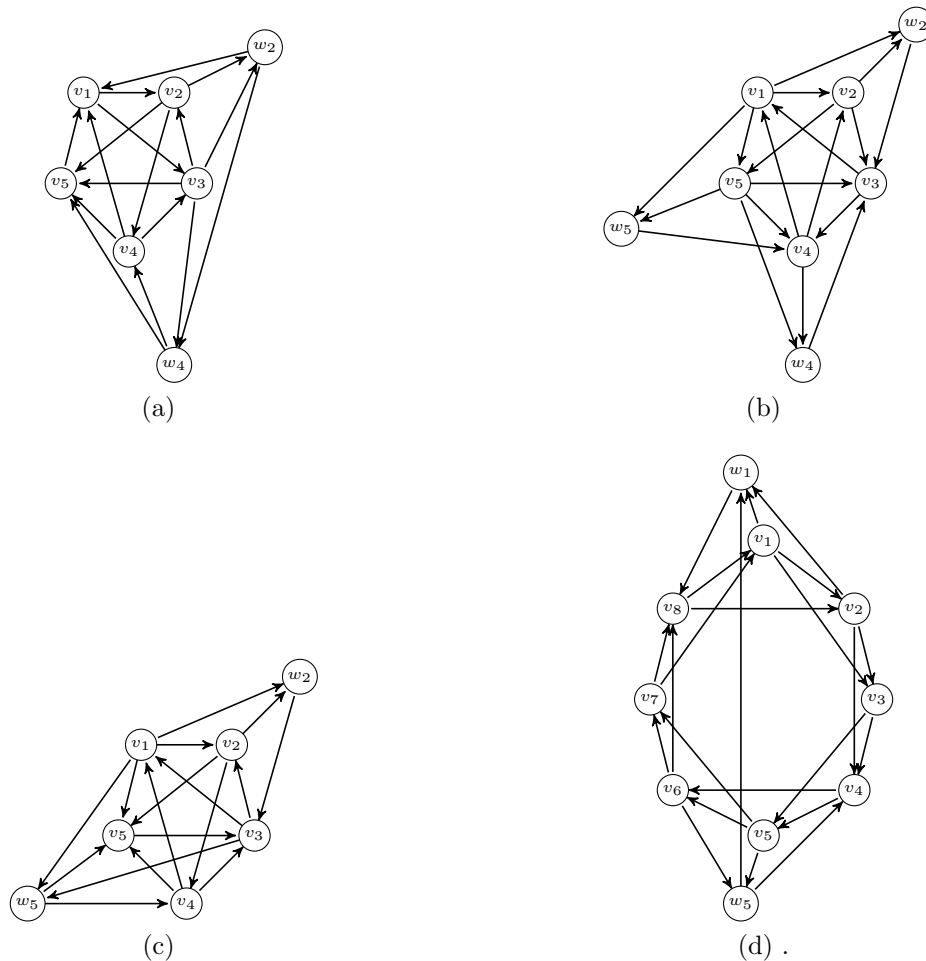


Figure 3: (a) The orientation for Lemma 13 has $|EE| = 57$ and $|EO| = 59$.
 (b) The orientation for Lemma 14 has $|EE| = 105$ and $|EO| = 106$.
 (c) The orientation for Lemma 15 has $|EE| = 63$ and $|EO| = 64$.
 (d) The orientation for Lemma 16 has $|EE| = 299$ and $|EO| = 298$.

program. (Ours is available at <https://dl.dropboxusercontent.com/u/8609833/Web/WebGraphs/WebGraphsTestPage.html>.) For the curious reader, in the appendix we provide hand checkable proofs of these lemmas.

Lemma 13. *Let H be a 5-cycle v_1, \dots, v_5 with pendant edges at v_2 and v_4 , leading to vertices w_2 and w_4 , respectively, and let w_2 and w_4 have a common neighbor x (off the cycle). Let $G = H^2 - x$; now G is d_1 -paintable.*

Proof. The orientation \vec{D} shown in Figure 3(a) has $|EE(\vec{D})| = 57$ and $|EO(\vec{D})| = 59$. Hence, the Alon–Tarsi Theorem implies that the graph is d_1 -paintable. \square

Lemma 14. *Let H be a 5-cycle v_1, \dots, v_5 with pendant edges at v_2 , v_4 , and v_5 , leading to vertices w_2 , w_4 , and w_5 , respectively. Let $G = H^2$; now G is d_1 -paintable.*

Proof. The orientation \vec{D} shown in Figure 3(b) has $|EE(\vec{D})| = 105$ and $|EO(\vec{D})| = 106$. Hence, the Alon–Tarsi Theorem implies that the graph is d_1 -paintable. \square

Lemma 15. *Let H be a 5-cycle v_1, \dots, v_5 with pendant edges at v_2 and v_5 , leading to vertices w_2 and w_5 , respectively, and let w_5 and v_3 have a common neighbor x (off the cycle). Let $G = H^2 - x$; now G is d_1 -paintable.*

Proof. The orientation \vec{D} shown in Figure 3(c) has $|EE(\vec{D})| = 63$ and $|EO(\vec{D})| = 64$. Hence, the Alon–Tarsi Theorem implies that the graph is d_1 -paintable. \square

Lemma 16. *8-cycle + two pendant edges + extra edge: Let J_8 consist of an 8-cycle on vertices v_1, \dots, v_8 (in clockwise order) with pendant edges at v_1 and v_5 leading to vertices w_1 and w_5 . Form \vec{D}_8 by squaring J_8 , adding the edge w_1w_5 and orienting the edges as follows. Orient edges $v_i v_{i+1}$ and $v_i v_{i+2}$ away from v_i (with subscripts modulo 8). Orient $w_1 v_8$ away from w_1 and $v_1 w_1$ and $v_2 w_1$ toward w_1 ; similarly, orient $w_5 v_4$ away from w_5 and $v_5 w_5$ and $v_6 w_5$ toward w_5 . Finally, orient $w_5 w_1$ toward w_1 . We will show that $\text{diff}(\vec{D}_8) \neq 0$ (or else $\text{diff}(\vec{D}_8 \setminus B) \neq 0$ for some subset $B \subseteq \{w_1, w_5\}$).*

Proof. The orientation \vec{D} shown in Figure 3(d) has $|EE(\vec{D})| = 299$ and $|EO(\vec{D})| = 298$. Hence, the Alon–Tarsi Theorem implies that the graph is d_1 -paintable. \square

Form \vec{P}_n from $(P_n)^2$ by orienting all edges from left to right. Number the vertices as v_1, \dots, v_n from left to right. A subgraph $\vec{T} \subseteq \vec{P}_n$ is *weakly eulerian* if each vertex $w \notin \{v_1, v_n\}$ satisfies $d^+(w) = d^-(w)$ and $d^+(v_1) = d^-(v_n) = i$ for some $i \in \{1, 2\}$. Let $EE_i(\vec{P}_n)$ (resp. $EO_i(\vec{P}_n)$) denote the set of even (resp. odd) weakly eulerian subgraphs where $d^+(v_1) = d^-(v_n) = i$. Finally, let $f_i(n) = |EE_i(\vec{P}_n)| - |EO_i(\vec{P}_n)|$. We will not apply the following lemma directly to find d_1 -paintable subgraphs. However, it will be helpful in the proofs for some of the remaining d_1 -paintable graphs, specifically those that include cycles of arbitrary length.

Lemma 17. *If $n = 3k + j$ for some positive integer k and $j \in \{-1, 0, 1\}$, then $f_1(n) = j$ and for $n \geq 4$ also $f_2(n) = -f_1(n - 2)$, with $f_i(n)$ as defined above.*

Proof. Rather than directly counting weakly eulerian subgraphs, we again use a parity-reversing bijection. We first prove that $f_2(n) = -f_1(n - 2)$. The complement of each $\vec{D} \in EE_2(\vec{P}_n) \cup EO_2(\vec{P}_n)$ has $d^+(v_2) = d^-(v_{n-1}) = 1$ and $d^+(w) = d^-(w)$ for each $w \notin \{v_1, v_2, v_{n-1}, v_n\}$ (and $d^+(v_1) = d^-(v_n) = d^-(v_2) = d^+(v_{n-1}) = 0$). Since \vec{P}_n has $2n - 3$ edges, each digraph has parity opposite its complement; so $f_2(n) = -f_1(n - 2)$.

Now we determine $f_1(n)$. Let \vec{T} be a weakly eulerian subgraph with $d^+(v_1) = 1$. Consider the directed paths $v_1 v_3$ and $v_1 v_2, v_2 v_3$. If \vec{T} contains all of one path and none of the other, then we can pair \vec{T} with its complement, which has opposite parity. If neither of these cases holds, then we must have $v_1 v_2, v_2 v_4 \in \vec{T}$ and $v_1 v_3, v_2 v_3 \notin \vec{T}$. This yields $f_1(n) = f_1(n - 3)$. It remains only to check that $f_1(2) = -1$, $f_1(3) = 0$, and $f_1(4) = 1$. \square

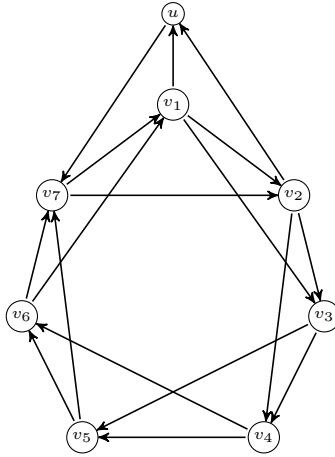


Figure 4: The orientation for Lemma 18 with $n = 7$.

Lemma 18. *Cycle + one pendant edge:* Let J_n consist of an n -cycle on vertices v_1, \dots, v_n (in clockwise order) with a pendant edge at v_1 leading to vertex u . Form \overrightarrow{D}_n by squaring J_n and orienting the edges as follows. Orient edges $v_i v_{i+1}$ and $v_i v_{i+2}$ away from v_i (with subscripts modulo n). Orient uv_n away from u and $v_1 u$ and $v_2 u$ toward u . We will show that $\text{diff}(\overrightarrow{D}_n) \neq 0$ when $n \not\equiv 2 \pmod{3}$ (or else $\text{diff}(\overrightarrow{D}_n - u) \neq 0$).

Proof. Form \overrightarrow{D}_n as in the lemma. We will show that $\text{diff}(\overrightarrow{D}_n) \neq 0$, and thus J_n^2 is d_1 -paintable. We may assume that $\text{diff}(\overrightarrow{D}_n - u) = 0$, for otherwise $\overrightarrow{D}_n - u$ is d_1 -paintable. Thus, restricting our count to the set A of circulations with $d^+(u) = 1$ does not affect the difference. Let \overrightarrow{T} be a circulation in A . Consider the directed paths $v_1 u$ and $v_1 v_2, v_2 u$. If \overrightarrow{T} contains all edges of one path and none of the other, then we can pair \overrightarrow{T} via a parity-reversing bijection. So we assume we are not in one of those cases. Clearly \overrightarrow{T} contains $\overrightarrow{uv_n}$ and exactly one of $v_1 u$ and $v_2 u$. Thus either (i) $v_2 u \in \overrightarrow{T}$ and $v_1 u, v_1 v_2 \notin \overrightarrow{T}$ or (ii) $v_1 u, v_1 v_2 \in \overrightarrow{T}$ and $v_2 u \notin \overrightarrow{T}$.

Case (i): $v_2 u \in \overrightarrow{T}$ and $v_1 u, v_1 v_2 \notin \overrightarrow{T}$. Since $v_2 u \in \overrightarrow{T}$ and $v_1 v_2 \notin \overrightarrow{T}$, we must have $v_n v_2 \in \overrightarrow{T}$ and $v_2 v_3, v_2 v_4 \notin \overrightarrow{T}$. By removing edges $uv_n, v_n v_2, v_2 u$, we see that these circulations are in bijection with the circulations in $\overrightarrow{D}_n - u - v_2$ (with the parity of each subgraph reversed). If we exclude the empty graph, these circulations are in bijection with those counted by $f_1(n-1)$, since $d^+(v_1) = 1$ and $d^-(v_3) = 1$. Adding 1 for the empty subgraph, this difference is $1 - f_1(n-1)$, and when we account for removing edges $uv_n, v_n v_2, v_2 u$, the difference is $-1 + f_1(n-1)$.

Case (ii): $v_1 u, v_1 v_2 \in \overrightarrow{T}$ and $v_2 u \notin \overrightarrow{T}$. Since $v_1 u, v_1 v_2 \in \overrightarrow{T}$, we have $v_{n-1} v_1, v_n v_1 \in \overrightarrow{T}$ and $v_1 v_3 \notin \overrightarrow{T}$. After removing edges $v_n v_1, v_1 u, uv_n$, we see that these circulations are in bijection with the circulations in $\overrightarrow{D}_n - u - v_n v_1 - v_1 v_3$ that contain edges $v_{n-1} v_1$ and $v_1 v_2$. We will count the difference of these even and odd circulations, then multiply the total by -1 (to account for removing edges $v_1 u, uv_n, v_n v_1$) before adding to the total above.

We consider two subcases: $v_n v_2 \notin \vec{T}$ and $v_n v_2 \in \vec{T}$. In the first case, these circulations are in bijection with circulations of $\vec{D}_{n-1} - u - v_1$ (since $d^+(v_n) = 0$ and v_1 may be suppressed). This difference is counted by $f_1(n-2)$. In the second case, the difference is counted by $-f_2(n)$, since we may think of deleting $v_1 v_2$ and replacing $v_n v_2$ with $v_n v_1$; our path now starts at v_2 and runs through v_n to v_1 (and the parity is changed when accounting for $v_1 v_2$).

Thus, the total difference in Case (ii) is counted by $f_1(n-2) - f_2(n)$. Thus, the total difference overall is counted by $-1 + f_1(n-1) - f_1(n-2) + f_2(n) = -1 + f_1(n-1) - 2f_1(n-2)$. Substituting values from Lemma 17 shows that this expression is non-zero when $n \not\equiv 2 \pmod 3$. \square

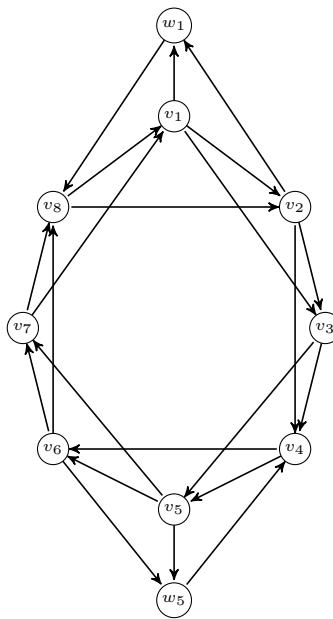


Figure 5: The orientation for Lemma 19 with $n = 8$.

Lemma 19. *Cycle + two pendant edges:* For $n \geq 7$, let J_n consist of an n -cycle on vertices v_1, \dots, v_n (in clockwise order) with pendant edges at v_1 and v_5 leading to vertices w_1 and w_5 . Form \vec{D}_n by squaring J_n and orienting the edges as follows. Orient edges $v_i v_{i+1}$ and $v_i v_{i+2}$ away from v_i (with subscripts modulo n). Orient $w_1 v_n$ away from w_1 and $v_1 w_1$ and $v_2 w_1$ toward w_1 ; similarly, orient $w_5 v_4$ away from w_5 and $v_5 w_5$ and $v_6 w_5$ toward w_5 . We will show that $\text{diff}(\vec{D}_n) \neq 0$ (or else $\text{diff}(\vec{D}_n \setminus B) \neq 0$ for some subset $B \subseteq \{w_1, w_5\}$).

Proof. Form \vec{D}_n as in the lemma. We will show that $\text{diff}(\vec{D}_n) \neq 0$, and thus J_n^2 is d_1 -paintable. For each nonempty $B \subseteq \{w_1, w_5\}$, we may assume that $\text{diff}(\vec{D}_n \setminus B) = 0$, for otherwise $\vec{D}_n \setminus B$ is d_1 -paintable. Thus, restricting our count to the set A of circulations with $d^+(w_1) = 1$ and $d^+(w_5) = 1$ does not affect the difference.

Let \vec{T} be a circulation in A . Clearly \vec{T} contains $\overrightarrow{w_1v_n}$ and exactly one of $\overrightarrow{v_1w_1}$ and $\overrightarrow{v_2w_1}$. Consider the directed paths v_1w_1 and v_1v_2, v_2w_1 . If \vec{T} contains all edges of one path and none of the other, then we can pair \vec{T} via a parity-reversing bijection. So we assume we are not in one of those cases. Thus either (i) $v_2w_1 \in \vec{T}$ and $v_1w_1, v_1v_2 \notin \vec{T}$ or (ii) $v_1w_1, v_1v_2 \in \vec{T}$ and $v_2w_1 \notin \vec{T}$.

Now we consider the directed paths v_5w_5 and v_5v_6, v_6w_5 . Among those circulations, within Cases (i) and (ii), where \vec{T} contains all of one path and none of the other we again pair \vec{T} via a parity-reversing bijection, by removing the edges of one path and adding the edges of the other. Thus, we need only consider two subcases in each case: (1) $v_6w_5 \in \vec{T}$ and $v_5w_5, v_5v_6 \notin \vec{T}$ and (2) $v_5w_5, v_5v_6 \in \vec{T}$ and $v_6w_5 \notin \vec{T}$.

Case (i.1): $v_2w_1 \in \vec{T}$ and $v_1w_1, v_1v_2 \notin \vec{T}$ and also $v_6w_5 \in \vec{T}$ and $v_5w_5, v_5v_6 \notin \vec{T}$. Since $v_2w_1 \in \vec{T}$, we must have $v_nv_2 \in \vec{T}$ and also $v_2v_3, v_2v_4 \notin \vec{T}$. Similarly, since $v_6w_5 \in \vec{T}$, we must have $v_4v_6 \in \vec{T}$ and also $v_6v_7, v_6v_8 \notin \vec{T}$. Since both triangles $w_1v_nv_2$ and $v_4v_6w_5$ must be included in every circulation under consideration, we may remove w_1, v_2, w_5, v_6 without changing the total difference. Now any non-empty circulation must contain both v_1v_3 and v_5v_7 . But we have a parity reversing bijection between those circulations containing v_3v_5 and those containing v_3v_4, v_4v_5 , so for non-empty circulations the difference is zero. Thus after adding in the empty circulation, we see that the total difference is 1 for this case.

Case (i.2): $v_2w_1 \in \vec{T}$ and $v_1w_1, v_1v_2 \notin \vec{T}$ and also $v_5w_5, v_5v_6 \in \vec{T}$ and $v_6w_5 \notin \vec{T}$. Since $v_2w_1 \in \vec{T}$, we must have $v_nv_2 \in \vec{T}$ and hence $v_2v_3, v_2v_4 \notin \vec{T}$. Since the triangle $w_1v_nv_2$ must be included in every circulation under consideration, we may remove w_1, v_2 at the cost of negating the difference. Since $v_5w_5, v_5v_6 \in \vec{T}$, we must have $w_5v_4, v_3v_5, v_4v_5 \in \vec{T}$ and $v_5v_7 \notin \vec{T}$. But then $v_3v_4 \notin \vec{T}$ and hence $v_4v_6 \notin \vec{T}$. Now we may remove w_5 and v_4 at the cost of negating the difference again. Now removing v_3 and v_5 we lose three edges that must be in every circulation and the resulting difference is counted by $f_1(n-4)$; the paths run from v_6 through v_n to v_1 . Hence this case contributes $-f_1(n-4)$ to the difference.

Case (ii.1): $v_1w_1, v_1v_2 \in \vec{T}$ and $v_2w_1 \notin \vec{T}$ and also $v_6w_5 \in \vec{T}$ and $v_5w_5, v_5v_6 \notin \vec{T}$. Since $v_1w_1, v_1v_2 \in \vec{T}$, we get $v_nv_1, v_{n-1}v_1 \in \vec{T}$. Since $v_6w_5 \in \vec{T}$ and $v_5v_6 \notin \vec{T}$, we get $v_4v_5 \in \vec{T}$ and $v_6v_7, v_6v_8 \notin \vec{T}$. Since we have $v_{n-1}v_1 \in \vec{T}$, we must also have $v_5v_7 \in \vec{T}$. Since $v_6v_7, v_6v_8 \notin \vec{T}$ and $v_5v_7 \in \vec{T}$, we get $d^+(v_2) = 1$. This also implies $d^+(v_{n-1}) = 1$. Now when $n \geq 9$ our difference is counted by $-f_1(3)f_1(n-7)$. Here $f_1(3)$ accounts for the edges of the path from v_2 to v_5 and $f_1(n-7)$ accounts for the edges of the path from v_7 to v_{n-1} (and the -1 accounts for the 9 edges that are present but not on either of these paths). Since $f_1(3) = 1$, the total for this case is $-f_1(n-7)$. When $n = 8$ the total is $-f_1(3) = -1$ and when $n = 7$ the total is 0, since $v_{n-1} = v_6$. Now by Lemma 17, together with checking the cases $n = 7$ and $n = 8$, we get that this case is counted by $-f_1(n-4)$.

Case (ii.2): $v_1w_1, v_1v_2 \in \vec{T}$ and $v_2w_1 \notin \vec{T}$ and also $v_5w_5, v_5v_6 \in \vec{T}$ and $v_6w_5 \notin \vec{T}$. Since $v_1w_1, v_1v_2 \in \vec{T}$, we must have $w_1v_n, v_nv_1, v_{n-1}v_1 \in \vec{T}$ and $v_1v_3 \notin \vec{T}$. Since $v_5w_5, v_5v_6 \in \vec{T}$, we must have $w_5v_4, v_3v_5, v_4v_5 \in \vec{T}$ and $v_5v_7 \notin \vec{T}$. Suppose $v_nv_2 \notin \vec{T}$.

Now $v_2v_4 \notin \vec{T}$, so $d^+(v_4) = 1$. Now our problem reduces to computing $-f_1(n - 6)$; the $f(n - 6)$ accounts for the edges on the path from v_6 to v_{n-1} and the -1 accounts for the 11 other edges that are present. Suppose instead that $v_nv_2 \in \vec{T}$. Now our problem reduces to computing $f_2(n - 4)$, accounting for the edges on the two paths from v_1 (after replacing v_nv_2 by v_nv_1) and the 12 edges present but not on these paths.

So, combining the contributions from all cases we get that the difference is $1 - f_1(n - 4) - f_1(n - 4) - f_1(n - 6) + f_2(n - 4)$. By Lemma 17 this is $1 - 2(f_1(n - 4) + f_1(n - 6)) \neq 0$ when $n \geq 8$. When $n = 7$ the difference is $1 - 2f_1(3) - 1 + f_2(3) = -1$. \square

For $n \geq 4$, a subgraph $\vec{T} \subseteq \vec{P}_n$ is *extra weakly eulerian* if each $w \notin \{v_1, v_2, v_{n-1}, v_n\}$ satisfies $d^+(w) = d^-(w)$, $d^+(v_1) = d^-(v_n) = 1$, $d^+(v_2) = d^-(v_2) + 1$ and $d^-(v_{n-1}) = d^+(v_{n-1}) + 1$. Let $EE^*(\vec{P}_n)$ (resp. $EO^*(\vec{P}_n)$) denote the set of even (resp. odd) extra weakly eulerian subgraphs. Finally, let $g(n) = |EE^*(\vec{P}_n)| - |EO^*(\vec{P}_n)|$. Lemma 20 is analogous to Lemma 17, but for extra weakly eulerian subgraphs.

Lemma 20. *If $n = 3k + j \geq 4$ for a positive integer k and $j \in \{-1, 0, 1\}$, then $g(n) = -j$.*

Proof. Let $\vec{T} \subseteq \vec{P}_n$ be extra weakly eulerian. Consider the directed paths v_1v_3 and v_1v_2, v_2v_3 . If \vec{T} contains all of one path but none of the other, then we can pair \vec{T} with its complement which has opposite parity. If neither of these cases holds, then we must have either $v_1v_3, v_2v_3 \in \vec{T}$ and $v_1v_2 \notin \vec{T}$ or $v_1v_2 \in \vec{T}$ and $v_1v_3, v_2v_3 \notin \vec{T}$. The latter case is impossible, so suppose we have $v_1v_3, v_2v_3 \in \vec{T}$ and $v_1v_2 \notin \vec{T}$. Then $v_3v_4, v_3v_5 \in \vec{T}$ and $v_2v_4 \notin \vec{T}$. Hence the difference is counted by $g(n - 3)$. It remains only to check that $g(4) = -1$, $g(5) = 1$ and $g(6) = 0$. \square

5 Generalizing to Alon–Tarsi number

Excepting the direct proofs of paintability in Section 4.1, we have actually proved that all the excluded subgraphs have a good Alon–Tarsi orientation. This suggests that the main theorem might hold more generally for the Alon–Tarsi number $AT(G)$ —the least k for which G has an orientation \vec{D} with $\Delta^+(\vec{D}) \leq k - 1$ and $EE(\vec{D}) \neq EO(\vec{D})$. Here we show that this is indeed the case.

Main Theorem for AT. *If G is a connected graph with maximum degree $\Delta \geq 3$ and G is not the Peterson graph, the Hoffman–Singleton graph, or a Moore graph with $\Delta = 57$, then $AT(G^2) \leq \Delta^2 - 1$.*

The proof is identical to the paintability proof except that we need to replace all the auxiliary lemmas with their AT counterparts. We first consider the two subgraph lemmas, which are actually easier to prove in the AT context.

Lemma 21. *Let G be a graph with maximum degree Δ and H be an induced subgraph of G that is d_1 -AT. If $G \setminus H$ is $(\Delta - 1)$ -AT, then G is $(\Delta - 1)$ -AT.*

Proof. Let G and H satisfy the hypotheses. Take an orientation of $G \setminus H$ demonstrating that it is $(\Delta - 1)$ -AT and an orientation of H demonstrating that it is d_1 -AT. Orient all edges between H and $G \setminus H$ into $G \setminus H$. Call the resulting oriented graph \vec{D} . Now \vec{D} satisfies the outdegree requirements to be $(\Delta - 1)$ -AT, since the outdegree of each vertex in $G \setminus H$ is unchanged and the outdegree of each $v \in V(H)$ has increased by $d_G(v) - d_H(v)$. Since no directed cycle in D has vertices in both H and $\vec{D} \setminus H$, the circulations of \vec{D} are just all pairings of circulations of H and $\vec{D} \setminus H$. Therefore $EE(\vec{D}) - EO(\vec{D}) = EE(H)EE(\vec{D} \setminus H) + EO(H)EO(\vec{D} \setminus H) - (EE(H)EO(\vec{D} \setminus H) + EO(H)EE(\vec{D} \setminus H)) = (EE(H) - EO(H))(EE(\vec{D} \setminus H) - EO(\vec{D} \setminus H)) \neq 0$. Hence G is $(\Delta - 1)$ -AT. \square

Lemma 22. *Let G be a graph with maximum degree Δ and let H be an induced subgraph of G^2 . If H is d_1 -AT, then G^2 is d_1 -AT. If there exists v with $d_{G^2}(v) < \Delta^2 - 1$, then G^2 is $(\Delta^2 - 1)$ -AT.*

Proof. We prove the first statement first. Form G' from G by contracting $V(H)$ to a single vertex r . Let T be a spanning tree in G' rooted at r . Let σ be an ordering of the vertices of $G \setminus H$ by nonincreasing distance in T from r . Take an orientation of H demonstrating that it is d_1 -AT; direct all edges between H and $G \setminus H$ towards $G \setminus H$ and direct all other edges of G^2 toward the vertex that comes earlier in σ . Call the resulting oriented graph \vec{D} . By construction, all circulations in \vec{D} are contained in H , so $EE(\vec{D}) \neq EO(\vec{D})$. Clearly, every vertex in \vec{D} has indegree at least two and hence G^2 is d_1 -AT.

Now we prove the second statement, which has a similar proof. Suppose there exists v with $d_{G^2}(v) < \Delta^2 - 1$. As before we order the vertices by nonincreasing distance in some spanning tree T from v , and we put v and some neighbor u last in σ . Since $d_{G^2}(v) < \Delta^2 - 1$, either (i) v lies on a 3-cycle or 4-cycle or else (ii) $d_G(v) < \Delta$ or v has some neighbor u with $d_G(u) < \Delta$; in Case (ii), by symmetry we assume $d_G(v) < \Delta$. In Case (i), $d_{G^2}(u) \leq \Delta^2 - 1$ for some neighbor u of v on the short cycle and by assumption $d_{G^2}(v) < \Delta^2 - 1$; so the two final vertices of σ are u and v . In Case (ii), we again have $d_{G^2}(v) < \Delta^2 - 1$ and $d_{G^2}(u) \leq \Delta^2 - 1$, so again u and v are last in σ . \square

The proof of Lemma 22 proves something slightly more general.

Corollary 23. *Let G be a graph with maximum degree Δ and let H be an induced subgraph of G^2 . Let $f(v) = d(v) - 1$ for each high vertex of G^2 and $f(v) = d(v)$ for each low vertex. If H is f -AT, then G^2 is $(\Delta^2 - 1)$ -AT.*

Now each of Lemmas 13–19 was already proved for AT. It remains to prove the lemmas in Section 4.1 for AT. We do this by exhibiting in Figure 6 a good Alon–Tarsi orientation for each. For brevity, we do not prove that the counts differ; instead we give the actual even/odd circulation counts for the reader to check at her leisure. Each vertex is labeled with its indegree for easy checking. Note that three of the cases in Lemma 7 are handled by Lemmas 10, 11, and 12 (none of which depend on Lemma 7). We conclude by generalizing the conjectures we mentioned in the introduction to the Alon–Tarsi number.

Conjecture 6 (Borodin-Kostochka Conjecture (Alon–Tarsi version)). *If G is a graph with $\Delta \geq 9$ and $\omega \leq \Delta - 1$, then $AT(G) \leq \Delta - 1$.*

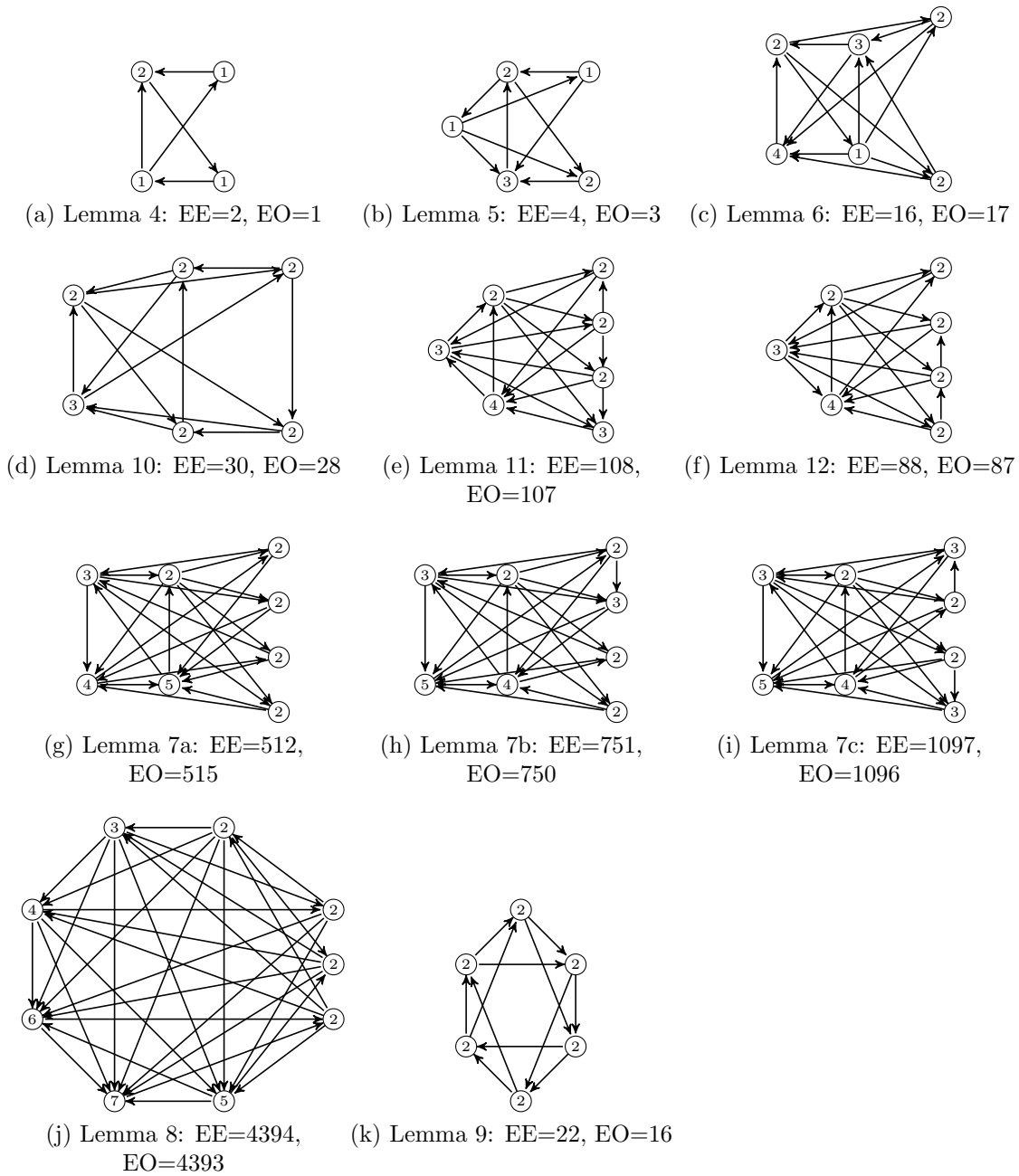


Figure 6: Good orientations for the AT versions of Lemmas 4–12.

A Proofs of Reducibility, without Computers

In Section 4.2, our proofs of Lemmas 13, 14, 15, and 16 consisted of simply providing orientations for which $|EE|$ and $|EO|$ differ. Here we give hand-checkable proofs.

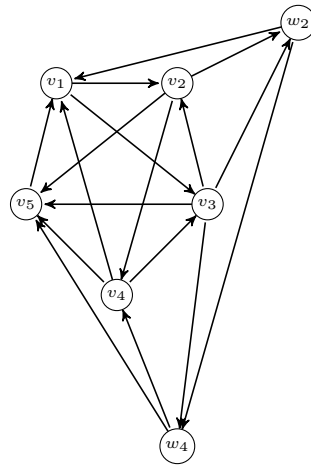


Figure 7: The orientation for Lemma 13.

Lemma 13. *Let H be a 5-cycle v_1, \dots, v_5 with pendant edges at v_2 and v_4 , leading to vertices w_2 and w_4 , respectively, and let w_2 and w_4 have a common neighbor x (off the cycle). Let $G = H^2 - x$; now G is d_1 -paintable.*

Proof. Orient G to form \vec{D} with these out-neighborhoods: $N^+(v_1) = \{v_2, v_3\}$, $N^+(w_2) = \{v_1, w_4\}$, $N^+(v_2) = \{w_2, v_4, v_5\}$, $N^+(v_3) = \{v_2, w_2, w_4, v_5\}$, $N^+(w_4) = \{v_4, v_5\}$, $N^+(v_4) = \{v_1, v_3, v_5\}$, $N^+(v_5) = \{v_1\}$. See Figure 7.

We will show that $\text{diff}(\vec{D}) \neq 0$. Since each vertex has at least two in-edges, this proves that G is d_1 -paintable. Let $R = \{v_3w_2, v_3w_4\}$. For any nonempty subset S of R , we must have $\text{diff}(\vec{D} \setminus R) = 0$. This is because each vertex on the 5-cycle has outdegree at most 3, so will get a list of size at most 4. And clearly, we cannot always color K_5 from lists of size at most 4. Thus, it suffices to count the difference, when restricted to the set A of circulations \vec{T} such that $v_3w_2, v_3w_4 \in \vec{T}$.

Let \vec{T} be such a circulation. Note that $v_3v_2, v_3v_5 \notin \vec{T}$, and thus $v_1v_3, v_4v_3 \in \vec{T}$. Now we consider the 8 possible subsets of $\{w_4v_4, w_4v_5, v_4v_5\}$ in \vec{T} . Clearly $d^+(w_4) \geq 1$ and $d^-(v_5) \leq 1$. Also, we can pair the case $w_4v_4, v_4v_5 \in \vec{T}$ and $w_4v_5 \notin \vec{T}$ with the case coming from its complement. Thus, we can restrict to the case when $w_4v_4 \in \vec{T}$ and $v_4v_5 \notin \vec{T}$ (and we're not specifying whether w_4v_5 is in or out). Now consider the directed triangle v_1v_2, v_2v_4, v_4v_1 . We can pair the cases when all or none of these edges are in \vec{T} . Thus we may assume that either exactly 1 or exactly 2 of these edges are in. Considering indegree and outdegree of v_2 shows that we must have $v_1v_2 \in \vec{T}$ and $v_2v_4, v_4v_1 \notin \vec{T}$. This implies $w_2v_1, v_5v_1 \in \vec{T}$. Now we have two ways to complete \vec{T} . We

can have $v_2w_2, w_2w_4, w_4v_5 \in \vec{T}$ and $v_2v_4 \notin \vec{T}$ or vice versa. Each of these gives $|E(\vec{T})|$ odd; thus, we get $|\text{diff}(D)| = 2$. \square

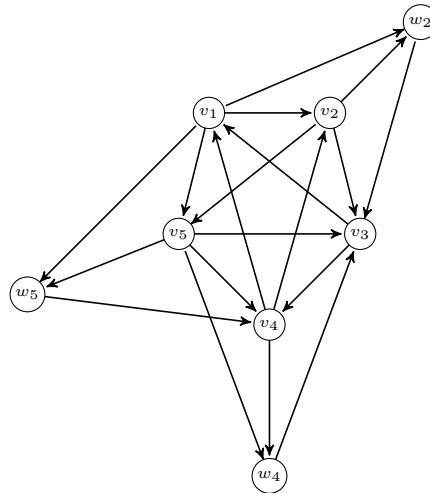


Figure 8: The orientation for Lemma 14.

Lemma 14. *Let H be a 5-cycle v_1, \dots, v_5 with pendant edges at v_2, v_4 , and v_5 , leading to vertices w_2, w_4 , and w_5 , respectively. Let $G = H^2$; now G is d_1 -paintable.*

Proof. We orient G to form \vec{D} with these out-neighborhoods: $N^+(v_1) = \{v_2, w_2, v_5, w_5\}$, $N^+(v_2) = \{w_2, v_3, v_5\}$, $N^+(w_2) = \{v_3\}$, $N^+(v_3) = \{v_1, v_4\}$, $N^+(v_4) = \{v_1, v_2, w_4\}$, $N^+(w_4) = \{v_3\}$, $N^+(v_5) = \{v_3, v_4, w_4, w_5\}$, $N^+(w_5) = \{v_4\}$. See Figure 8.

We will show that $\text{diff}(\vec{D}) \neq 0$. Since each vertex has at least two in-edges, this proves that G is d_1 -paintable. If $\text{diff}(\vec{D} - w_2) \neq 0$, then we are done, since $\vec{D} - w_2$ is d_1 -paintable. Thus, we can assume that $\text{diff}(\vec{D} - w_2) = 0$. Similarly, we can assume that $\text{diff}(\vec{D} \setminus S) = 0$ for every $S \subseteq \{w_2, w_4, w_5\}$. Thus, it suffices to count the difference, when restricted to the set A of circulations such that $d^+(w_2) = 1$, $d^+(w_4) = 1$, and $d^+(w_5) = 1$. Let \vec{T} be such a circulation. So $w_2v_3, w_4v_3, w_5v_4 \in \vec{T}$. Now $d^+(v_3) = 2$, so $v_3v_1, v_3v_4 \in \vec{T}$ and $v_2v_3, v_5v_3 \notin \vec{T}$. In particular, $d^-(v_1) \geq 1$, so $d^+(v_1) \geq 1$.

Now we will pair some circulations in A via a parity-reversing bijection. Consider the paths v_1w_2 and v_1v_2, v_2w_2 . If a circulation contains all edges in one path and none in the other, then we can pair it via a bijection. The same is true for the paths v_1w_5 and v_1v_5, v_5w_5 . Since $1 \leq d^+(v_1) \leq 2$, and also $d^-(w_2) = d^-(w_5) = 1$, the only way that \vec{T} can avoid these cases is if either (i) $v_1v_2, v_1w_2 \in \vec{T}$ or (ii) $v_1v_5, v_1w_5 \in \vec{T}$. Before we consider these cases, note that in each case $v_4v_1 \in \vec{T}$.

Case (i): Now we must have $v_1w_5, v_1v_5 \notin \vec{T}$. Note that $v_2w_2 \notin \vec{T}$, which implies $v_4v_2 \notin \vec{T}$. Also $v_2v_5 \in \vec{T}$. Further, $d^-(w_5) = 1$ implies $v_5w_5 \in \vec{T}$, which in turn yields

$v_5v_4, v_5w_4 \notin \vec{T}$. Finally, $v_4w_4 \in \vec{T}$. Thus, we have a unique \vec{T} (with an odd number of edges).

Case (ii): Now we must have $v_1w_2, v_1v_2 \notin \vec{T}$ and also $v_5w_5 \notin \vec{T}$. Note that $v_2w_2 \in \vec{T}$, which implies that $v_4v_2 \in \vec{T}$ and also that $v_2v_5 \notin \vec{T}$. Now we get that either (a) $v_5v_4 \in \vec{T}$, and thus $v_4w_4 \in \vec{T}$ and $v_5w_4 \notin \vec{T}$ or else (b) $v_5w_4 \in \vec{T}$ and $v_5v_4, v_4w_4 \notin \vec{T}$. Again, by a parity-reversing bijection, we see that together these circulations contribute nothing to $\text{diff}(A)$ (in fact there is only one of each). Now combining Cases (i) and (ii), we get that $|\text{diff}(A)| = 1$, and in fact $|\text{diff}(\vec{D})| = 1$. Thus, G is d_1 -paintable. \square

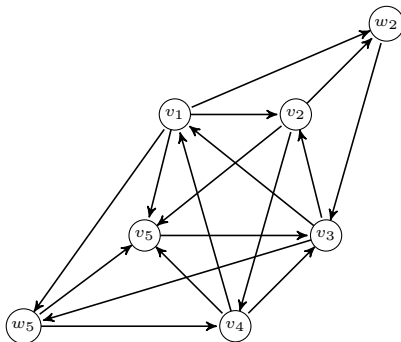


Figure 9: The orientation for Lemma 15.

Lemma 15. *Let H be a 5-cycle v_1, \dots, v_5 with pendant edges at v_2 and v_5 , leading to vertices w_2 and w_5 , respectively, and let w_5 and v_3 have a common neighbor x (off the cycle). Let $G = H^2 - x$; now G is d_1 -paintable.*

Proof. We orient G to form \vec{D} with these out-neighborhoods: $N^+(v_1) = \{v_2, w_2, v_5, w_5\}$, $N^+(v_2) = \{w_2, v_4, v_5\}$, $N^+(w_2) = \{v_3\}$, $N^+(v_3) = \{v_1, v_2, w_5\}$, $N^+(v_4) = \{v_1, v_3, v_5\}$, $N^+(v_5) = \{v_3\}$, $N^+(w_5) = \{v_4, v_5\}$. See Figure 9.

We will show that $\text{diff}(\vec{D}) \neq 0$. Since each vertex has at least two in-edges, this proves that G is d_1 -paintable. Note that for each nonempty subset $S \subseteq \{w_2, w_5\}$, we have $\text{diff}(\vec{D} \setminus S) = 0$, since otherwise we can color the corresponding subgraph from lists of size 4, even though it contains a 5-clique. So by inclusion-exclusion, we can restrict our count of diff to the set of circulations A where w_2 and w_5 each have positive indegree. Consider the paths v_1w_2 and v_1v_2, v_2w_2 . Let \vec{T} be a circulation in A . If \vec{T} contains all edges of one path and none of the other, then we can pair it via a parity-reversing bijection. So we assume we are not in these situations. Since w_2 has positive indegree, and hence indegree 1, we either have (i) $v_1w_2, v_1v_2 \in \vec{T}$ and $v_2w_2 \notin \vec{T}$ or (ii) $v_2w_2 \in \vec{T}$ and $v_1w_2, v_1v_2 \notin \vec{T}$.

Case (i): $v_1w_2, v_1v_2 \in \vec{T}$ and $v_2w_2 \notin \vec{T}$. Clearly $w_2v_3 \in \vec{T}$. Since $d^+(v_1) = 2$, we have $v_3v_1, v_4v_1 \in \vec{T}$ and $v_1v_5, v_1w_5 \notin \vec{T}$. Suppose $v_3v_2 \in \vec{T}$. Now also $v_2v_4, v_2v_5, v_5v_3 \in \vec{T}$. Finally, since w_5 has positive indegree, $v_3w_5, w_5v_4, v_4v_3 \in \vec{T}$. The resulting circulation is

even. Suppose instead that $v_3v_2 \notin \vec{T}$. If $v_2v_5 \in \vec{T}$, then we get $v_5v_3, v_3w_5, w_5v_4 \in \vec{T}$. The resulting circulation is odd. If instead $v_2v_5 \notin \vec{T}$ and $v_2v_4 \in \vec{T}$, then we have three possibilities to ensure $d^+(w_5) > 0$. Either $v_3w_5, w_5v_4, v_4v_5, v_5v_3 \in \vec{T}$ or $v_3w_5, w_5v_4, v_4v_3 \in \vec{T}$ or $v_3w_5, w_5v_5, v_5v_3 \in \vec{T}$. Two of the resulting circulations are odd and one is even. Thus in total for Case (i), we have one more odd circulation than even.

Case (ii): $v_2w_2 \in \vec{T}$ and $v_1w_2, v_1v_2 \notin \vec{T}$. We have $v_2w_2 \in \vec{T}$, which implies $w_2v_3 \in \vec{T}$ and $v_3v_2 \in \vec{T}$. This further yields $v_2v_4, v_2v_5 \notin \vec{T}$. Again we will pair some of the circulations in A via a parity-reversing bijection. Consider the paths v_3w_5 and v_3v_1, v_1w_5 . If a circulation contains all edges in one path and none in the other, then we can pair it via a bijection. Since $1 \leq d^-(w_5)$, the only way that \vec{T} can avoid these cases is if either (a) $v_1w_5 \in \vec{T}$ and $v_3v_1 \notin \vec{T}$ or (b) $v_3v_1 \in \vec{T}$ and $v_1w_5 \notin \vec{T}$ (and thus $v_3w_5 \in \vec{T}$ or (c) $v_3v_1, v_1w_5, v_3w_5 \in \vec{T}$. Consider (a). $v_1w_5 \in \vec{T}$ implies $v_4v_1 \in \vec{T}$, and thus $w_5v_4 \in \vec{T}$. We also have the option of all or none of v_3w_5, w_5v_5, v_5v_3 in \vec{T} . One of the resulting circulations is odd and the other is even. Consider (b). Now $v_3v_1 \in \vec{T}$ and $v_1w_5 \notin \vec{T}$ imply $v_1v_5 \in \vec{T}$, and thus $v_5v_3 \in \vec{T}$. Now $d^+(w_5) > 0$ implies $v_3w_5, w_5v_4, v_4v_3 \in \vec{T}$. The resulting circulation is odd. Consider (c). Now we get $w_5v_5 \in \vec{T}$, which implies $v_5v_3 \in \vec{T}$. We also get $w_5v_4 \in \vec{T}$, which implies $v_4v_3 \in \vec{T}$. The resulting circulation is even. Thus in total for Case (ii), we have the same number of even and odd circulations.

So combining Cases (i) and (ii), we have one more odd circulation than even. Thus $\text{diff}(\vec{D}) \neq 0$, so G is d_1 -paintable. \square

Form \vec{P}_n from $(P_n)^2$ by orienting all edges from left to right. Number the vertices as v_1, \dots, v_n from left to right. A subgraph $\vec{T} \subseteq \vec{P}_n$ is *weakly eulerian* if each vertex $w \notin \{v_1, v_n\}$ satisfies $d^+(w) = d^-(w)$ and $d^+(v_1) = d^-(v_n) = i$ for some $i \in \{1, 2\}$. Let $EE_i(\vec{P}_n)$ (resp. $EO_i(\vec{P}_n)$) denote the set of even (resp. odd) weakly eulerian subgraphs where $d^+(v_1) = d^-(v_n) = i$. Finally, let $f_i(n) = |EE_i(\vec{P}_n)| - |EO_i(\vec{P}_n)|$. We will not apply the following lemma directly to find d_1 -paintable subgraphs. However, it will be helpful in the proof for the remaining d_1 -paintable graph, which includes cycles of arbitrary length.

Lemma 16. *8-cycle + two pendant edges + extra edge: Let J_8 consist of an 8-cycle on vertices v_1, \dots, v_8 (in clockwise order) with pendant edges at v_1 and v_5 leading to vertices w_1 and w_5 . Form \vec{D}_8 by squaring J_8 , adding the edge w_1w_5 and orienting the edges as follows. Orient edges $v_i v_{i+1}$ and $v_i v_{i+2}$ away from v_i (with subscripts modulo 8). Orient w_1v_8 away from w_1 and v_1w_1 and v_2w_1 toward w_1 ; similarly, orient w_5v_4 away from w_5 and v_5w_5 and v_6w_5 toward w_5 . Finally, orient w_5w_1 toward w_1 . We will show that $f(\vec{D}_8) \neq 0$ (or else $f(\vec{D}_8 \setminus B) \neq 0$ for some subset $B \subseteq \{w_1, w_5\}$).*

Proof. Form \vec{D}_8 as in the lemma. Suppose $f(\vec{D}_8 \setminus B) = 0$ for each subset $\emptyset \neq B \subseteq \{w_1, w_5\}$. Then by Lemma 19, we have $\text{diff}(\vec{D}_8 - w_5w_1) \neq 0$. Hence it will suffice to show that the circulations of \vec{D}_8 containing w_5w_1 are half odd and half even.

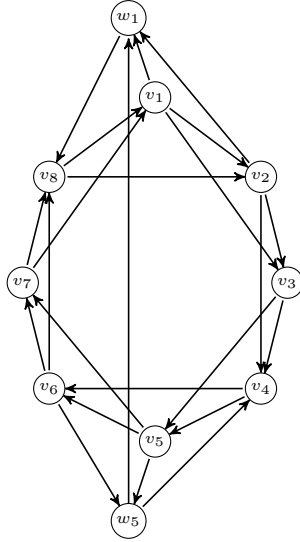


Figure 10: The orientation for Lemma 16.

Let \vec{T} be a circulation of \vec{D}_8 containing w_5w_1 . Then $w_1v_8 \in \vec{T}$ and $v_1w_1, v_2w_1 \notin \vec{T}$. After suppressing w_1 , we are looking at all circulations containing w_5v_8 .

Consider the directed paths v_5w_5 and v_5v_6, v_6w_5 . If \vec{T} contains all edges of one path and none of the other, then we can pair \vec{T} via a parity-reversing bijection. So we assume we are not in one of those cases. Thus either (i) $v_6w_5 \in \vec{T}$ and $v_5w_5, v_5v_6 \notin \vec{T}$, (ii) $v_5w_5, v_5v_6 \in \vec{T}$ and $v_6w_5 \notin \vec{T}$, (iii) $v_5w_5, v_5v_6, v_6w_5 \in \vec{T}$ or (iv) $v_6w_5, v_5w_5 \in \vec{T}$ and $v_5v_6 \notin \vec{T}$.

Case (i): $v_6w_5 \in \vec{T}$ and $v_5w_5, v_5v_6 \notin \vec{T}$. Then $v_4v_6 \in \vec{T}$ and $w_5v_4, v_6v_7, v_6v_8 \notin \vec{T}$. Now we can suppress v_6 and w_5 . First suppose $v_5v_7 \notin \vec{T}$. Now $v_7, v_5 \notin \vec{T}$ and what remains is counted by $-f_1(5)$. Instead suppose $v_5v_7 \in \vec{T}$. Then the difference is counted by $g(7)$; the path is from v_7 to v_5 . Hence the total difference is $g(7) - f_1(5) = -1 - (-1) = 0$.

Case (ii): $v_5w_5, v_5v_6 \in \vec{T}$ and $v_6w_5 \notin \vec{T}$. Then $v_3v_5, v_4v_5 \in \vec{T}$ and $w_5v_4, v_5v_7 \notin \vec{T}$. Now we can suppress w_5 . First suppose $v_4v_6 \in \vec{T}$. There is only one possible circulation and it contains all edges except v_7v_8 ; this circulation is odd, hence the difference is -1 . Now suppose $v_4v_6 \notin \vec{T}$. If $v_6v_7 \in \vec{T}$, then $v_6v_8 \notin \vec{T}$ and the difference is counted by $-g(6)$; the path is from v_7 to v_4 . If $v_6v_7 \notin \vec{T}$, then $v_6v_8, v_8v_1, v_8v_2 \in \vec{T}$ and $v_7 \notin \vec{T}$. Now the difference is counted by $-g(4)$; the path is from v_1 to v_4 . Hence the total difference is $-1 - g(6) - g(4) = 0$.

Case (iii): $v_5w_5, v_5v_6, v_6w_5 \in \vec{T}$. Then $w_5v_4, v_3v_5, v_4v_5 \in \vec{T}$ and $v_5v_7 \notin \vec{T}$. If $v_4v_6, v_6v_7 \in \vec{T}$, then the difference is counted by $g(6)$; the path is from v_7 to v_4 . Since $v_6v_7 \in \vec{T}$ and $v_4v_6 \notin \vec{T}$ is impossible, we may assume either $v_4v_6 \in \vec{T}$ and $v_6v_7 \notin \vec{T}$ or $v_4v_6, v_6v_7 \notin \vec{T}$. Suppose we are in the former case. Then $v_6v_8, v_8v_1, v_8v_2 \in \vec{T}$ and $v_7 \notin \vec{T}$.

This difference is counted by $g(4)$; the path is from v_1 to v_4 . Now suppose $v_4v_6, v_6v_7 \notin \vec{T}$. Then $v_7 \notin \vec{T}$ and $v_6v_8 \notin \vec{T}$. This difference is counted by $f_1(4)$; the path is from v_8 to v_3 . Hence the total difference is $g(6) + g(4) + f_1(4) = 0$.

Case (iv): $v_6w_5, v_5w_5 \in \vec{T}$ and $v_5v_6 \notin \vec{T}$. Then $w_5v_4, v_4v_6 \in \vec{T}$ and $v_6v_7, v_6v_8 \notin \vec{T}$. If $v_5v_7 \notin \vec{T}$, then $v_7 \notin \vec{T}$ and the difference is counted by $f_1(6) = 0$; the path is from v_8 to v_5 . Hence we may assume $v_5v_7 \in \vec{T}$. Then $v_3v_5, v_4v_5 \in \vec{T}$ and the difference is counted by $g(6) = 0$; the path is from v_7 to v_4 .

So in each of the four cases, half the circulations are even and half are odd. Thus, the difference is not affected by the circulations that use edge w_5w_1 . Now by Lemma 19, $f(\vec{D}) \neq 0$, so \vec{D} is d_1 -paintable. \square

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