

# Improving Failure Responsiveness in Jini Leasing

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## Abstract

*Distributed systems require strategies to detect and recover from failures. Many protocols for distributed systems employ a strategy based on leases, which grant a leaseholder access to data or services for a limited time (the lease period). Choosing an appropriate lease period involves tradeoffs among resource utilization, responsiveness, and system size. We explain these tradeoffs for Jini Network Technology. Then, we describe an adaptive algorithm that enables a Jini system, given a fixed allocation of resources, to vary lease periods with system size to achieve the best responsiveness. We anticipate that similar procedures could improve failure responsiveness in other distributed systems that rely on leases. We describe how we implemented our adaptive algorithm in “reggie”, a publicly available implementation of the Jini lookup service. We can use our implementation to demonstrate how adaptive leasing provides the best available responsiveness as network size varies.*

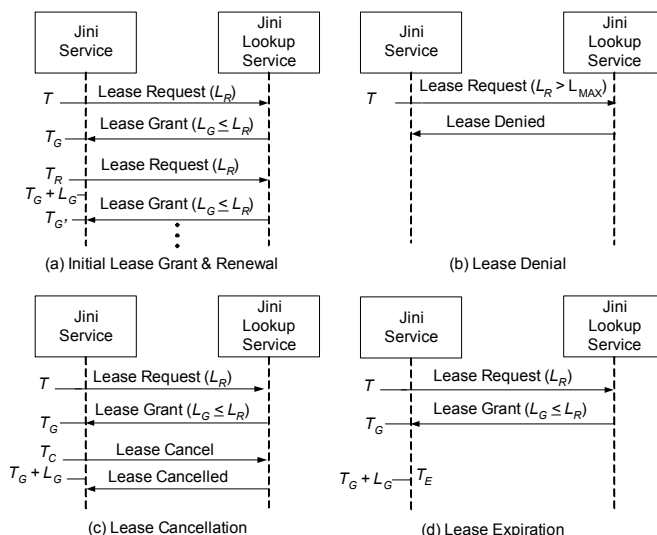
## 1. Introduction

Distributed systems require strategies to detect and recover from failures. One commonly used strategy employs a leasing mechanism, where a node grants a leaseholder access to a resource for a limited time (the lease period). Once the resource is no longer needed, the leaseholder may relinquish its lease. If the resource is needed beyond the original lease period, then the leaseholder can renew the lease by requesting additional lease periods. If the leaseholder does not renew before expiration of the lease period, the lease grantor assumes leaseholder failure and terminates the lease.

Choosing an appropriate lease period entails tradeoffs among resource utilization, responsiveness, and number of leaseholders. We explore these issues in the context of service-discovery protocols, which allow distributed software components to discover each other and compose themselves into assemblies. Though several service-discovery protocols currently exist [e.g., 1-4], we selected Jini Network Technology [1] to demonstrate our ideas, because leasing plays a central role in registering Jini services. We base our analysis on the Jini specification [2].

## 2. Jini Leasing

Jini defines an architecture that enables clients and services to rendezvous through a third party, known as a lookup service. A Jini service registers a description of itself with each discovered lookup service. A Jini client may register a request to be notified by a lookup service of arriving or departing services of interest, or of changes in the attributes describing services of interest.



**Fig. 1. Message exchanges for four Jini leasing scenarios.**

Figure 1 illustrates message exchanges for some typical Jini leasing scenarios. A registering component requests registration for duration  $L_R$ , which may be accepted at time  $T_G$  for a granted lease period  $L_G \leq L_R$ .  $L_R$  may be any, which allows any value for  $L_G$ . To extend registration beyond  $L_G$ , registering components must renew the lease prior to an expiration time  $T_E = T_G + L_G$ ; otherwise, registration is revoked. This cycle continues until a Jini component cancels or fails to renew a lease. Lookup services assign  $L_G$  within a configured range,  $L_{MIN} \leq L_G \leq L_{MAX}$ . While a granted lease may not be revoked prior to  $T_E$ , lookup services may deny any lease request.

We can analyze performance of a Jini leasing system. Let  $S_R$  be lease-request size,  $S_G$  be lease-grant size, and  $N$  be

number of leaseholders. Typically, a leaseholder and lookup service exchange one request-grant pair per renewal cycle, with rate  $1/L_G$  Hz. Assuming identical  $L_G$  assigned for each lease, bandwidth use ( $B$ ) can be estimated as:  $B=(N/L_G)\cdot(S_R+S_G)$ . Assuming constant  $S_R$  and  $S_G$ ,  $B$  increases linearly with  $N$  and decreases exponentially with  $L_G$ . Another metric, responsiveness,  $R$ , measures the latency with which lookup services can detect leaseholder failure. Assuming uniformly distributed failure times, then expected responsiveness is  $R = L_G / 2$ ; thus,  $R$  is independent of  $N$ , but  $B$  and  $R$  are related through  $L_G$ .

These relationships can be used to constrain and predict behavior of a leasing system. For example, assume known requirements for  $R$  and  $B$ . The responsiveness equation can be rewritten to determine  $L_G$  [i.e.,  $L_G = 2R$ ]. Then, using  $L_G$ , the bandwidth equation can be transformed to find maximum system size [i.e.,  $N_{MAX} = (B \cdot L_G) / (S_R + S_G)$ ]. With this information, lookup services could grant lease periods  $\leq L_G$  to ensure required responsiveness, deny requested leases that would consume an excess share of bandwidth, and deny requests for leases once  $N$  reaches  $N_{MAX}$ .

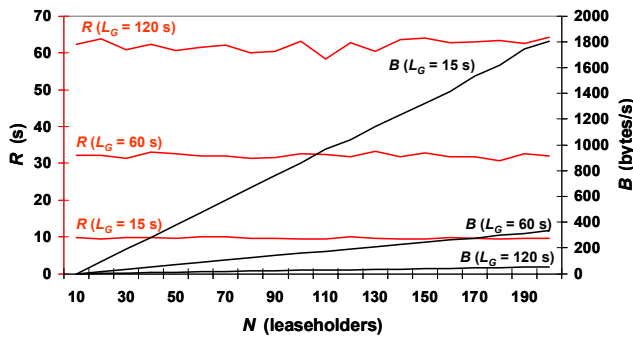


Fig. 2. System responsiveness ( $R$ ) – left-hand y-axis – and bandwidth usage ( $B$ ) – right-hand y-axis – for three granted lease periods ( $L_G = 15$  s,  $60$  s, and  $120$  s) as system size increases ( $N = 10$  to  $200$  leaseholders).

### 3. A Self-adaptive Algorithm for Jini Leasing

We propose an algorithm that restricts lease requests to  $L_R = any$ . Assuming a leasing system must consume at most bandwidth  $B$  and guarantee minimum average responsiveness  $R_{MIN}$ , a lookup service can grant a maximum lease period  $L_{MAX} = 2R_{MIN}$ . Given  $B$ ,  $S_R$ , and  $S_G$ , we can determine a maximum lease-renewal rate  $G = B / (S_R + S_G)$ . For minimum system size,  $N_{MIN} = 1$ , the lookup service can grant a minimum lease period  $L_{MIN} = 1/G$ . While this value for  $L_{MIN}$  respects the bandwidth constraint, other factors should be considered. For example, at  $L_{MIN} = 1/G$  leaseholder processing burden might prove unacceptable. Instead, a leasing system might constrain maximum responsiveness ( $R_{MAX}$ ), giving a minimum lease period  $L_{MIN} = 2R_{MAX}$ .

Knowing  $N$ , a lookup service may select a suitable granted lease period from a range ( $L_{MIN} \leq L_G \leq L_{MAX}$ ) using a simple algorithm. First, compute  $L_G = N/G$ . If  $L_G > L_{MAX}$ , then deny the lease; otherwise, if  $L_G < L_{MIN}$ , then set  $L_G = L_{MIN}$ . Assigning  $L_G$  with this algorithm permits a leasing system to constrain  $B$  and guarantee minimum average responsiveness ( $R_{MIN}$ ), while providing the best responsiveness achievable (up to  $R_{MAX}$ ) as  $N$  varies over  $1..N_{MAX}$ .

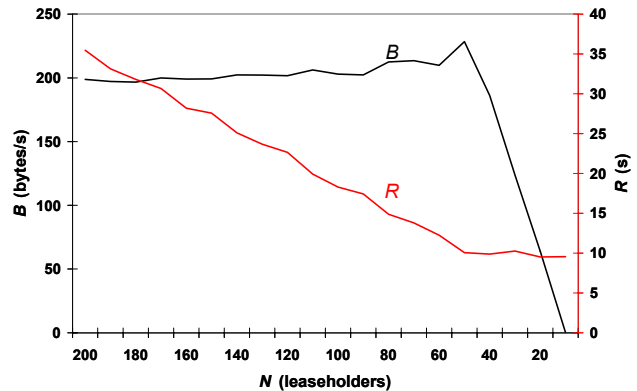


Fig. 3. Responsiveness ( $R$ ) – left-hand y-axis – and bandwidth usage ( $B$ ) – right-hand y-axis – as system size decreases ( $N = 200$  to  $0$  leaseholders).

### 4. Simulation Results

We coded an SLX discrete-event simulation [5] model of Jini to confirm our analysis and to investigate dynamic behavior of our self-adaptive algorithm. We conducted simulation experiments, varying  $N$  from  $10..200$  and  $L_G$  from  $15..300$  s in  $15$ -s increments. We used  $S_R = 128$  and  $S_G = 32$  bytes. Figure 2 shows simulated results for average  $B$  and  $R$  when  $L_G = 15$  s,  $60$  s, and  $120$  s. The simulation confirms our analyses: (1)  $B$  increases linearly with  $N$  for a given  $L_G$  and decreases exponentially with  $L_G$  for a given  $N$  and (2)  $R = L_G/2$ , independent of  $N$ . Next, we simulated our adaptive leasing algorithm. Figure 3 illustrates how the algorithm constrains  $B$  while improving  $R$  as system size decreases. These promising results led us to implement our adaptive algorithm in “reggie”, a publicly available implementation of a Jini lookup service.

### 5. Implementation in “reggie”

We base our adaptive-leasing implementation on the “reggie” lookup server provided with the Sun Microsystems Jini release. The “reggie” server implements the Jini specification for a lookup service, and includes additional extensions to allow remote administration of the lookup server through a service proxy. Administrative actions occur through the *RegistrarAdmin* interface, which is not part of the

Jini core specification, but a Sun extension to Jini (**com.sun.jini.reggie** package). The *RegistrarAdmin* interface allows basic monitoring, configuration, and control of an operational lookup server just as if it were any other type of Jini-enabled service. Using the *RegistrarAdmin* interface, an administrator can also perform some basic manipulation of minimum and maximum granted lease periods for services and events maintained on the lookup server. However, this method requires constant human supervision to optimize leasing performance in a Jini network. Such human supervision would prove impractical in a large network where numerous services may join and leave.

To implement adaptive leasing, we modified the “reggie” server implementation to assign lease grant times ( $L_G$ ) based on the required failure responsiveness ( $R$ ) of the system and the bandwidth ( $B$ ) allocated to lease renewal transactions. We added a collection of access methods to the *RegistrarAdmin* interface, callable via remote-method invocation (RMI), allowing a Jini client to view: the current  $L_{MIN}$ ,  $L_{MAX}$ , and  $L_G$ , the number of leaseholders ( $N$ ) on the server, the instantaneous average bandwidth ( $B_{AVG}$ ) consumed by lease renewals, and the instantaneous average failure responsiveness ( $R_{AVG}$ ). Since values for granted lease periods can be adjusted from changes to the allocated bandwidth and target responsiveness, we added methods to set  $B$  and  $R$  in the *RegistrarAdmin* interface.

The lookup server (**com.sun.jini.reggie.RegistrarImpl**) starts with default values for  $L_{MIN}$  and  $L_{MAX}$ . A Jini client can use the *RegistrarAdmin* interface to adjust target responsiveness and allocated bandwidth. Based on these adjustments, the lookup server computes new values for  $L_{MIN}$  and  $L_{MAX}$ . At regular intervals, the lookup service samples average bandwidth use ( $B_{AVG}$ ) and the number of leaseholders, adjusting granted lease periods ( $L_G$ ) accordingly. Current bandwidth usage is calculated by multiplying the number of lease transactions (RMI calls) by the size of messages involved in the transaction. Currently, the lookup service records these values once every sixty seconds of operation, or when an administrator changes  $R$  or  $B$ .

When a Jini service registers with the lookup service, it may either request a specific lease interval or use Jini’s *LEASE.ANY* constant to allow the lookup server to select an appropriate lease period for the service. In our implementation, if the service requests the *LEASE.ANY* constant, the lookup service uses the current value for  $L_G$  as the granted lease period. Otherwise, if the service requests a lease period in the range of  $L_{MIN}$  and  $L_{MAX}$ , it is granted. The lookup service rejects requests for leases outside this range.

Figure 4 shows a snapshot of a Jini client graphical user interface (GUI) that uses the modified “reggie” *RegistrarAdmin* interface. The left-hand column plots values for  $L_G$ , for bandwidth used ( $B_{AVG}$ ), and for average responsiveness ( $R_{AVG}$ ) over time. These graphs display values

returned to the client from regular polling of access methods *RegistrarAdmin* interface, which uses RMI to call the corresponding method in the lookup server. The right-hand column of Figure 4 lists leaseholders using the lookup service, and displays their current status in the Jini network. Note that the GUI displays only the current *known* status, as a service may have left the network, but a proxy could still be registered with the server. The proxy would be purged when its lease expires. The GUI in Figure 4 does not include the *RegistrarAdmin* GUI used to adjust allocated bandwidth or target responsiveness in the lookup service.

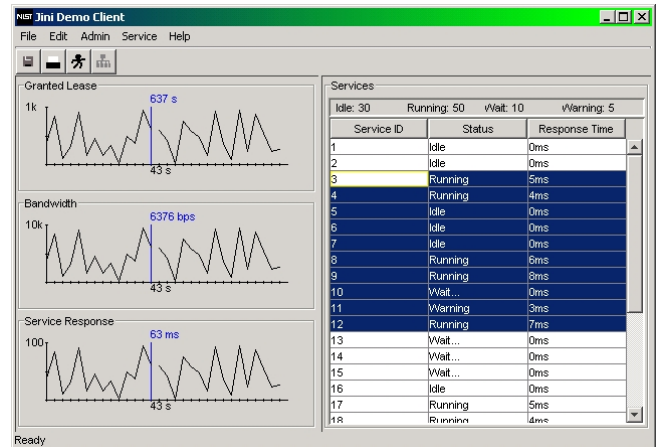


Fig. 4. Sample Graphical User Interface for a Jini client monitoring a lookup server that implements adaptive leasing

## 6. Acknowledgments

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## 7. References

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